

# Exploiting Software Vulnerabilities

## Software Vulnerabilities

### INTEGER OVERFLOWS AND FORMAT STRING BUGS

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**Master's Degree in Informatics Engineering**

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# Outline

- 1** Integer Vulnerabilities
  - Wraparound
  - Overflow
  - Truncation
  - Signedness errors
  
- 2** Format String Vulnerability
  - Formatted Output Functions
  - Description of the Vulnerability
  - Mitigation Strategies

# Outline

## 1 Integer Vulnerabilities

- Wraparound
- Overflow
- Truncation
- Signedness errors

## 2 Format String Vulnerability

# Integer Vulnerabilities

## Types of integer vulnerabilities

### ■ Overflow

- Occurs at runtime when the result of an integer expression exceeds the maximum value for its respective type
- For instance, two 8-bit unsigned integers may require up to 16 bits

### ■ Underflow

- Occurs at runtime when the result of an integer expression is less than its minimum value. So it wraps to the maximum integer for the integer type
- For instance, subtracting  $0 - 1$  and storing the result in a 16-bit unsigned integer will result in a value of  $2^{16} - 1$  (not  $-1$ )

### ■ Truncation

- Occurs when assigning an integer with a greater width to a smaller width
- For instance, converting an `int` to a `short` discards the leading bits of the `int` value, resulting in a (potential) information loss

### ■ Signedness error

- Occurs when a signed integer is interpreted as unsigned, or vice versa
- In the two-complement representation, such conversions cause the sign bit to be interpreted as the most significant bit (i.e.,  $2^{32} - 1 \neq -1$ )

# Integer Vulnerabilities

## Examples

```
struct pixmap {
    unsigned char *p;
    int x;
    /* xsize */
    int y;
    /* ysize */
    int bpp;
};
typedef struct pixmap pix;
.....
void readpgm(char *name, pix * p) {
    /* read pgm */...
    pnm_readpaminit(fp, &inпам);
    p->x=inпам.width;
    p->y=inпам.height;
    if(!(p->p=(char *)malloc(p->x*p->y)))
        F1("Error at malloc");
    for(i=0; i<inпам.height; i++){
        pnm_readpamrow(&inпам, tuplerow);
        for(j = 0; j<inпам.width; j++)
            p->p[i*inпам.width+j]=sample;
    }
}

void getComm(unsigned int len, char *src){
    unsigned int size;
    size = len - 2;
    char *comm = (char *)malloc(size + 1);
    memcpy(comm, src, size);
    return;
}
```

```
static inline u32 *decode_fh(u32 *p, struct svc_fh *fhp) {
    int size;
    fh_init(fhp, NFS3_FHSIZE);
    size = ntohl(*p++);
    if (size > NFS3_FHSIZE)
        return NULL;
    memcpy(&fhp->fh_handle.fh_base, p, size);
    fhp->fh_handle.fh_size = size;
    return p + XDR_QUADLEN(size);
}
```

```
int detect_attack(u_char *buf, int len, u_char *IV){
    static word16 *h = (word16 *) NULL;
    static word16 n = HASH_MIN_ENTRIES;
    register word32 i, j;
    word32 l;
    ...
    for(l=n; l<HASH_FACTOR(len/BSIZE); l=l<<2);
    if (h == NULL) {
        debug("Install crc attack detector.");
        n = l;
        h = (word16 *) xmalloc(n*sizeof(word16));
    } else
        for (c=buf, j=0; c<(buf+len); c+=BSIZE, j++){
            for (i = HASH(c) & (n - 1); h[i] != UNUSED; i = (i + 1) & (n - 1))
                ...;
            h[i] = j;
        }
}
```

# Integer Vulnerabilities

## Exploiting integer bugs

### Usually, they are indirectly exploited

#### ■ Arbitrary code execution

- For instance, insufficient memory allocation exploited by buffer overflows, heap overflows, overwrite attacks, etc.

#### ■ Denial of Service

- For instance, excessive memory allocation or infinite loops

#### ■ Attacks to bypass sanitization

- For instance, skipping an upper bounds check that ignores unexpected negative integer values

#### ■ Logic errors

- For instance, manipulating the reference counter by forcing a referenced object to be freed prematurely

**Further reading:** *Understanding Integer Overflow in C/C++*. W. Dietz et al. ACM Trans. Softw. Eng. Methodol. 25, 1, Article 2

(December 2015) doi: 10.1145/2743019

# Integer Vulnerabilities

## Consequences

### ■ **Silent break**

- Compiler optimizations can result in the exploitation of undefined behavior

### ■ **Time bombs**

- Today works, but improvements in optimization technologies can take advantage of it

### ■ **Illusion of predictability**

- Some compilers, at some optimization levels, have predictable behavior for some undefined operations

### ■ **Informal dialects**

- Some compilers have flags to force the two-complement behavior on signed overflows

### ■ **Non-standard standards**

- Meaning of overflow changes between standards (e.g.,  $1 \ll 31$ )
  - Implementation defined in ANSI C and C++98
  - Undefined by C99 and C11

# Integer vulnerabilities

## Wraparound – unsigned data types

- **Any calculation involving unsigned operands can never overflow**
- *What if the result of an operation cannot be represented by the resulting unsigned integer type?*
  - **The result is reduced by modulo the number that is one greater than the largest value that can be represented by the resulting type**
- Can occur with **addition and multiplication operations**
  - *n*-bit addition/subtraction operations require  $n + 1$  bits of precision
  - Similarly, multiplying *n*-bit requires  $2n$  bits of precision



# Integer vulnerabilities

## Wraparound – unsigned data types

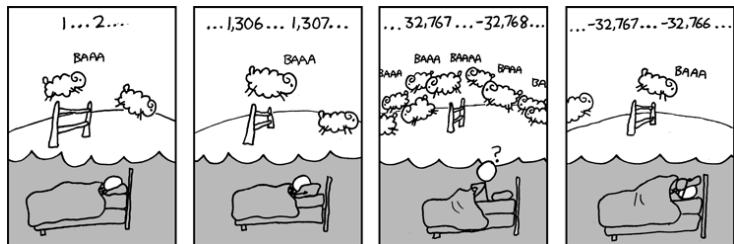
### How to avoid them?

- Check the wraparound, either before performing the operation that would cause it to occur or after
- `<limits.h>` limits are useful, but note that **naive use of them does not work:**

```
unsigned i, j, sum;
if (sum + i > UINT_MAX)
    // Too big error
else
    sum += i;
```

# Integer vulnerabilities

## Overflow – signed data types



- Occurs when a signed integer operation results in a value that cannot be represented in the resulting type
- **Signed integer overflow is undefined behavior in C, allowing implementations to silently wrap (the most common behavior), trap, or both**
- Since signed integer overflow produces a silent wraparound in most existing C compilers, **some programmers assume that this is a well-defined behavior**

# Integer vulnerabilities

## Find and fix overflows

### Shift operations

- **Operand values are checked for limits**
- If the verification passes, then the shift is made

### Arithmetic operations

- **Problem:  $n$ -bit addition/subtraction operations require  $n + 1$  bits of precision**
- Similarly, multiplying  $n$ -bit requires  $2n$  bits of precision
  - In C, addition, subtraction, multiplication, negation, and division **can result in overflow or underflow**
  - When  $-2^{n-1}$  is negated or divided by  $-1$ , the result overflows and wraps back to  $-2^{n-1}$
- **Three different methods** can be applied

# Integer vulnerabilities

## Find and fix overflows

### Detecting overflow for an operation on two signed integers $s_1$ and $s_2$

#### ■ Precondition test:

$$\begin{aligned} & ((s_1 > 0) \wedge (s_2 > 0) \wedge (s_1 > (\text{INT\_MAX} - s_2))) \vee \\ & ((s_1 < 0) \wedge (s_2 < 0) \wedge (s_1 < (\text{INT\_MAX} - s_2))) \end{aligned}$$

#### ■ Post-condition test of CPU flag (overflow flag)

#### ■ Post-condition test of width extension

- Convert  $s_1$  and  $s_2$  to a broader type
- Perform the operation
- Check if the result is within the limits w.r.t. the original (narrower) type

# Integer vulnerabilities

## Truncation

- Occurs as a result of an assignment or cast from a type with a larger width to a type with a smaller width
- Data may be lost if the value cannot be represented in the resulting type
  - For instance, adding `c1` and `c2` in the following program fragment produces a value outside the limits of `unsigned char`, considering an implementation where `unsigned char` is represented using 8 bits ( $2^8 - 1 = 255$ )

```
unsigned char sum, c1, c2;
```

```
c1 = 200;
```

```
c2 = 90;
```

```
sum = c1 + c2;
```

# Integer vulnerabilities

## Signedness errors

### Conversion rules

- **A set of rules that provides a mechanism to produce a common type** when
  - Both operands of a binary operator are balanced with a common type
  - The second and third arguments of the conditional operator ( ? : ) are balanced with a common type
- Balancing conversions involve two operands of different types
  - One or both operands can be converted
- **Many operators that accept integer operands perform conversions using the usual arithmetic conversions**, including \*, /, %, +, -, <, >, <=, >=, ==, !=, &, ^, |, and the condition operator ? :

# Integer vulnerabilities

## Signedness errors

### Example of conversion on x86-32

```
unsigned int ui = UINT_MAX;
signed char c = -1;

if (c == ui) {
    printf("-1 = 4,294,967,295?\n");
}
```

# Integer vulnerabilities

## Signedness errors

### Some notes

- **Implicit conversions simplify C programming**
- **Please note: conversions have the potential to lose or misinterpret data**
- Avoid conversions that result in:
  - **Loss of value** (cast to a type where the magnitude of the value cannot be represented)
  - **Loss of sign** (cast from signed type to unsigned type resulting in loss of sign)
- **Conversions of integers to a type with greater range and the same signedness are guaranteed safe for all data values in all compliant implementations**



# Outline

## 1 Integer Vulnerabilities

## 2 Format String Vulnerability

- Formatted Output Functions
- Description of the Vulnerability
- Mitigation Strategies

# Formatted output functions

- **Consist of a format string and a variable number of arguments**
  - The format string provides a set of instructions that are interpreted by the formatted output function
- **By controlling the content of the format string, a user can control execution of the formatted output function**

# Formatted output functions

- **Consist of a format string and a variable number of arguments**
  - The format string provides a set of instructions that are interpreted by the formatted output function
- By controlling the content of the format string, **a user can control execution of the formatted output function**

## Format string

- **Character sequences consisting of ordinary characters (excluding %) and conversion specifiers**
- Ordinary characters **are copied unchanged to the output stream**
- Conversion specifiers **indicate how to convert arguments and write the results to the output stream**
  - Starts with the %, interpreted from left to right

# Formatted output functions

## Conversion specifiers

`%[flags] [width] [.precision] [{length-modifier}] conversion-specifier`

- **Additional arguments are ignored**, if the number of arguments is greater than the conversion specifiers
- **Results are undefined**, if the number of arguments is less than the conversion specifiers

# Format string vulnerability

- Known since 1999 and exploited since 2000
- Unlike BOF, it is considered a programming error
- Easy to find
- Exploitation techniques are basic

# Format string vulnerability

## Real examples

Application	Found by	Impact	Years
wu-ftp 2.*	security.is	remote root	> 6
Linux rpc.statd	security.is	remote root	> 4
IRIX telnetd	LSD	remote root	> 8
Qualcomm Popper 2.53	security.is	remote user	> 3
Apache + PHP3	security.is	remote user	> 2
NLS / locale	CORE SDI	local root	?
screen	Jouko Pynnönen	local root	> 5
BSD chpass	TESO	local root	?
OpenBSD fstat	ktwo	local root	?

Adapted from <https://crypto.stanford.edu/cs155/papers/formatstring-1.2.pdf>

# Format string vulnerability

When does it appear?

**When user input is included in an ANSI C format function (in part or in full)**

# Format string vulnerability

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When user input is included in an ANSI C format function (in part or in full)

```
void error(char *s)
{
    fprintf(stderr, s);
}
```

*What if \*s is equal to "%s%s%s%s%s%s"?*



# Format string vulnerability

When does it appear?

When user input is included in an ANSI C format function (in part or in full)

```
void error(char *s)
{
    fprintf(stderr, s);
}
```

*What if \*s is equal to "%s%s%s%s%s%s"?*

- **Program will crash (most likely):** Denial-of-Service
- **Otherwise, the contents of the stack will be printed:** *privacy issues*

# Format string vulnerability

## Unsafe functions

- `fprintf`
- `printf`
- `sprintf`
- `snprintf`
- `vfprintf`
- `vprintf`
- `vsprintf`
- `vsnprintf`
- `setproctitle`
- `syslog`
- Others like `err*`, `verr*`, `warn*`, `vwarn*`

# Format string vulnerability

## ■ Functionality

- Simple conversion from C datatypes to string representation
- The representation format can be specified
- The resulting string is processed (e.g., output to stdout, stderr, syslog, etc.)

## ■ How does a format function work?

- **The format string, in fact, controls the behavior of the function**
  - Type of parameters to print
  - **Parameters are stored on the stack (pushed), either directly (value) or indirectly (reference)**
- 
- **The calling function knows how many parameters were pushed, as it has to make stack correction after returning** (by calling convention, more on this later)

# Format string vulnerability

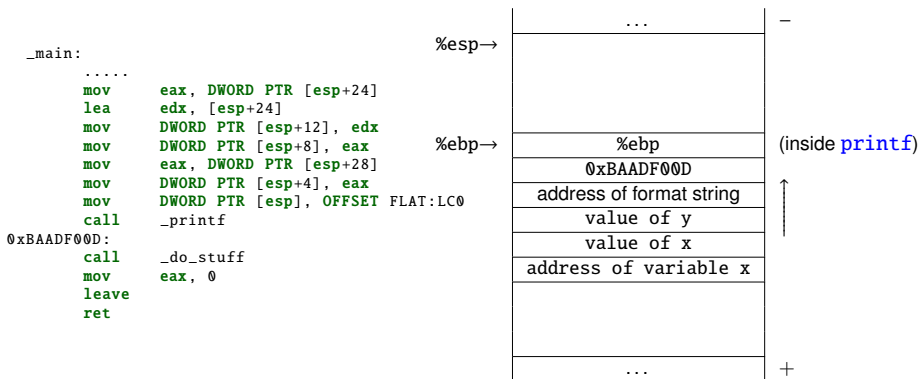
Specifier	Output	Example	Passed as
%d or i	Signed decimal integer	392	value
%u	Unsigned decimal integer	7235	value
%o	Unsigned octal	610	value
%x	Unsigned hexadecimal integer	7fa	value
%X	Unsigned hexadecimal integer (uppercase)	7FA	value
%f	Decimal floating point, lowercase	392.65	value
%F	Decimal floating point, uppercase	392.65	value
%e	Scientific notation (mantissa/exponent), lowercase	3.9265e+2	value
%E	Scientific notation (mantissa/exponent), uppercase	3.9265E+2	value
%g	Use the shortest representation: %e or %f	392.65	value
%G	Use the shortest representation: %E or %F	392.65	value
%a	Hexadecimal floating point, lowercase	-0xc.90fep-2	value
%A	Hexadecimal floating point, uppercase	-0XC.90FEP-2	value
%c	Character	a	value
%s	String of characters	sample	reference
%p	Pointer address	b8000000	value
%n	Nothing printed. The corresponding argument must be a pointer to a signed int. The number of characters written so far is stored in the pointed location.		reference
%%	A % followed by another % character will write a single % to the stream.	%	-

Taken from <http://www.cplusplus.com/reference/cstdio/printf/>

# Format string vulnerability

## Example

```
printf("Values %d, %d, %08x\n", y, x, &x);  
do_stuff(); // this call is in 0xBAADF00D address
```



# Format string vulnerability

**Channeling problem** *When two different types of information channels are merged into one and special escape characters (or sequences) are used to distinguish which channel is currently active*

- One channel is **data channel** (not parsed, just copied): **output strings**
- The other channel is a **control channel: format specifiers**

**NOTE:** *channeling issues are not security holes, but they make the bugs exploitable*

## Types

- **Type 1:** format string is partially user-supplied
- **Type 2:** a user-supplied string partially or completely is indirectly passed to a formatted output function

# Format string vulnerability

## Examples by type

### Type 1

```
char tmpbuf[512];
snprintf (tmpbuf, sizeof (tmpbuf), "foo: %s", user);
tmpbuf[sizeof (tmpbuf) - 1] = '\\0';
syslog (LOG_NOTICE, tmpbuf);
```

### Type 2

```
int Error (char *fmt, ...);
...
int someotherfunc (char *user)
{
    ...
    Error(user);
    ...
}
```

# Format string vulnerability

## Feasible attacks

### **Denial of Service**

- Reading an unallocated memory address crashes the application



# Format string vulnerability

## Feasible attacks

### Denial of Service

- Reading an unallocated memory address crashes the application

### Read-what-where

- Interesting format specifiers: %s, %p
  - Walk up reading the entire contents of the stack

```
printf("%s"); // will print the top of the stack  
printf("%5$s"); // will print the 5th element of the stack
```

# Format string vulnerability

## Feasible attacks

### Denial of Service

- Reading an unallocated memory address crashes the application

### Read-what-where

- Interesting format specifiers: `%s`, `%p`
  - Walk up reading the entire contents of the stack

```
printf("%s"); // will print the top of the stack  
printf("%5$s"); // will print the 5th element of the stack
```

### Write-what-where

- Interesting format specifier: `%n`
  - Writes to a specified variable **the number of bytes already written**

```
int i;  
  
printf("foobar%n", (int *)&i); // after, i = 6
```

# Format string vulnerability

## Exploitation goals

### ■ **Overwriting the Global Offset Table**

- Holds the address of each library function used by an ELF program
- Independent of stack or heap particulars
- **The control-flow execution is hijacked right after the program invokes the override function**

### ■ **DTORS** (in programs compiled with GNU GCC)

- Destructor table section (termination routines)
- Overwrite the pointer to the shellcode and thus **the control-flow execution is hijacked when the program exits**

### ■ **C library hooks**

- Present in the GNU C library and other proprietary libraries
- Hooks legitimately used by memory profiling and debugging tools
- **Control-flow execution is hijacked when the program invokes `malloc`, `realloc`, `realloc`, etc. (before the actual function is run)**

### ■ **`__atexit` structures**

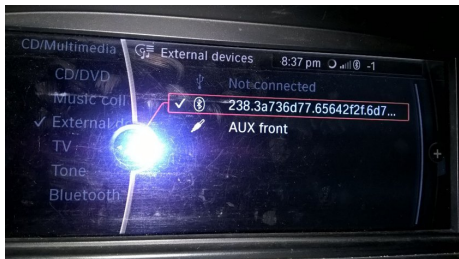
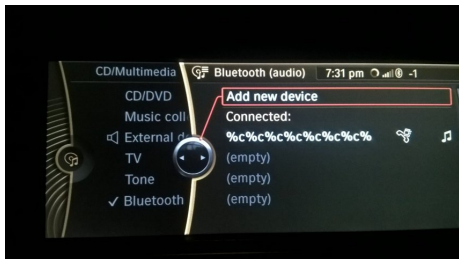
- Generic handler, **runs when a Linux program invokes `exit`**

### ■ **Function pointers**

- Commonly used by daemons for command processing or to simulate termination routine handlers

# Format string vulnerability

Real example today: 2011 BMW 330i



Associated CVE: CVE-2017-9212 (it also exists in Mercedes-Benz AMG, CVE-2020-16142)

Credits: [https://twitter.com/\\_\\_Obzy\\_\\_/status/864704956116254720](https://twitter.com/__Obzy__/status/864704956116254720)

# Format string vulnerability

Real example today: Audi A7 2014



**Credits:** <https://tiger-team-1337.blogspot.com/2020/10/audi-a7-2014-mmi-mishandles-format.html>

# Mitigation strategies

## Dynamic format strings

- Design your code so that **the user selects from a preexisting format string, rather than incorporating user input directly into the format string**

## Byte written restriction

- **Avoid buffer overflows that restrict the number of bytes written by formatted output functions**
- **Use the precision field as part of the %s conversion specifier**
  - `sprintf(buf, "Wrong command: %s\n", user);` ⇒  
`sprintf(buf, "Wrong command: %.495s\n", user);`
  - Even better if you use `snprintf` instead

# Mitigation strategies

## C11 Annex K functions

- Provide functions `fprintf_s`, `printf_s`, `snprintf_s`, `sprintf_s`, `vfprintf_s`, `vprintf_s`, `vsprintf_s`, `vsnprintf_s`
- Differ from their non-`_s` counterparts by:
  - **Not compatible with the `%n` format conversion specifier**
  - Constraint violation if pointers are null or format string is invalid
- **PLEASE NOTE:** these functions cannot avoid format string vulnerabilities that crash a program or are exploited to view memory
  - That is, **they only prevent from write-what-where attacks**

## GNU C compiler flags

- `-Wformat`, `-Wformat-nonliteral`, `-Wformat-security`
  - `-Wformat` checks calls to formatted output functions, examine the format string, and checks that the correct number and types of arguments are supplied
  - `-Wformat-nonliteral` warns if the format string is not a literal string and cannot be verified
  - `-Wformat-security` warns about calls to formatted output function that represent potential security issues

# Exploiting Software Vulnerabilities

## Software Vulnerabilities

### INTEGER OVERFLOWS AND FORMAT STRING BUGS

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