From UML State-Machine Diagrams to Erlang

Ricardo J. Rodríguez, Lars-Åke Fredlund, Ángel Herranz

All wrongs reversed

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Facultad de Informática, Universidad Complutense de Madrid

Outline

1. Motivation
2. A Transformation Approach
3. Related Work
4. Conclusions and Future Work

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Motivation (I)

Software Development Life-Cycle

- Phased involved for developing (and maintaining) software systems and code
- Deployment (production): what happens when faults (or changes) raise?

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- Model-Driven Engineering
  - Increase productivity, simplifying design
  - Maximise compatibility between systems

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Verify correctness BEFORE deployment

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Verify correctness BEFORE deployment

How?: Using model-checking

- Proofs of correctness
- Counter-examples (why not correct)

Motivation (II)

We have mixed...

- UML: standard *de facto* as modelling language
  - UML State Machines (UML-SMs): dynamic system behaviour
  - Assumption: intercommunication through asynchronous channels
- Erlang: functional and concurrent programming language
  - Native support for concurrency, distribution and fault tolerance
  - Concurrency based on asynchronous message passing
  - Widely used in the industry: T-Mobile, Ericsson, FB, WhatsApp...

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Contributions

- Minimise development time
  - Automatically generate Erlang code from UML-SMs

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Contributions

- Minimise development time
  - Automatically generate Erlang code from UML-SMs
- Detect problems in early stages (save efforts and costs)
  - Apply Erlang-based model checking techniques into UML-SMs

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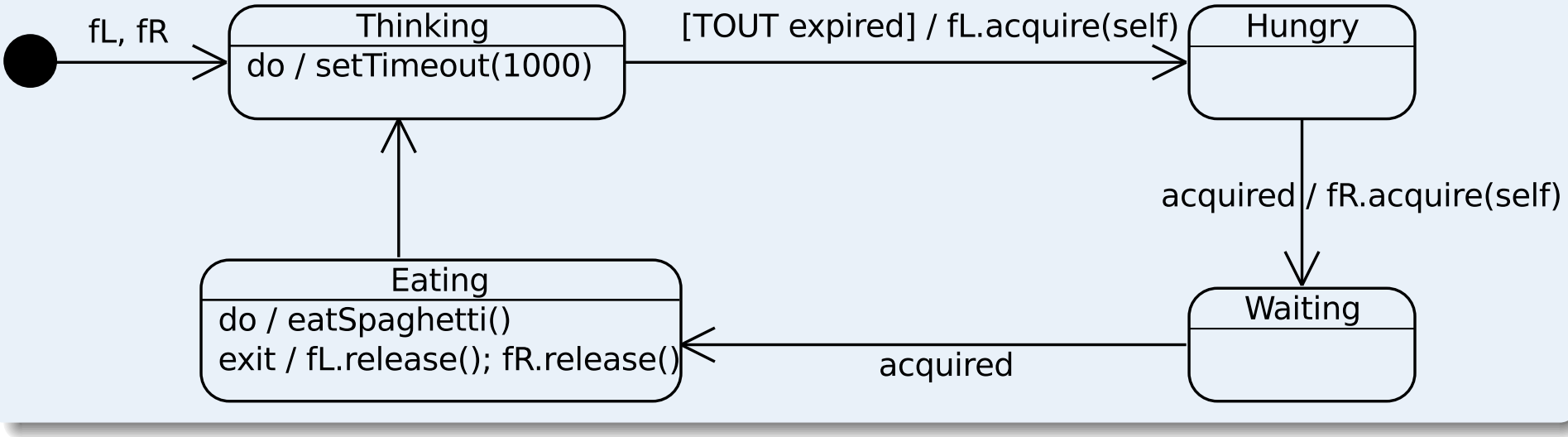
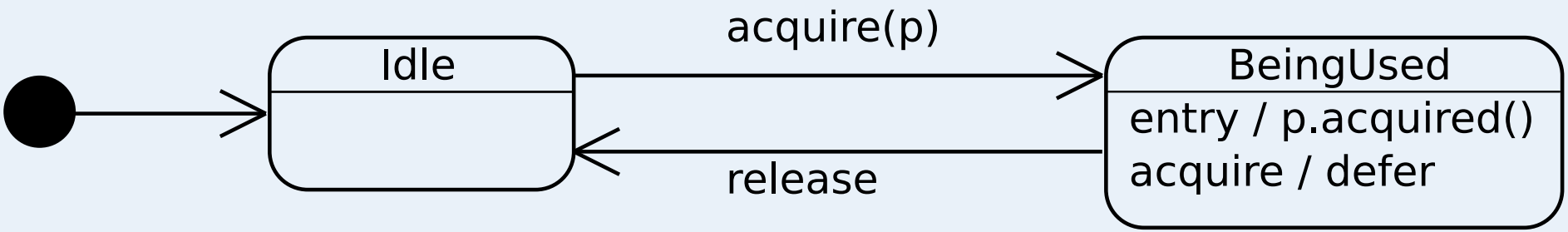
A Transformation Approach: The Dining Philosophers (I)

Philosopher



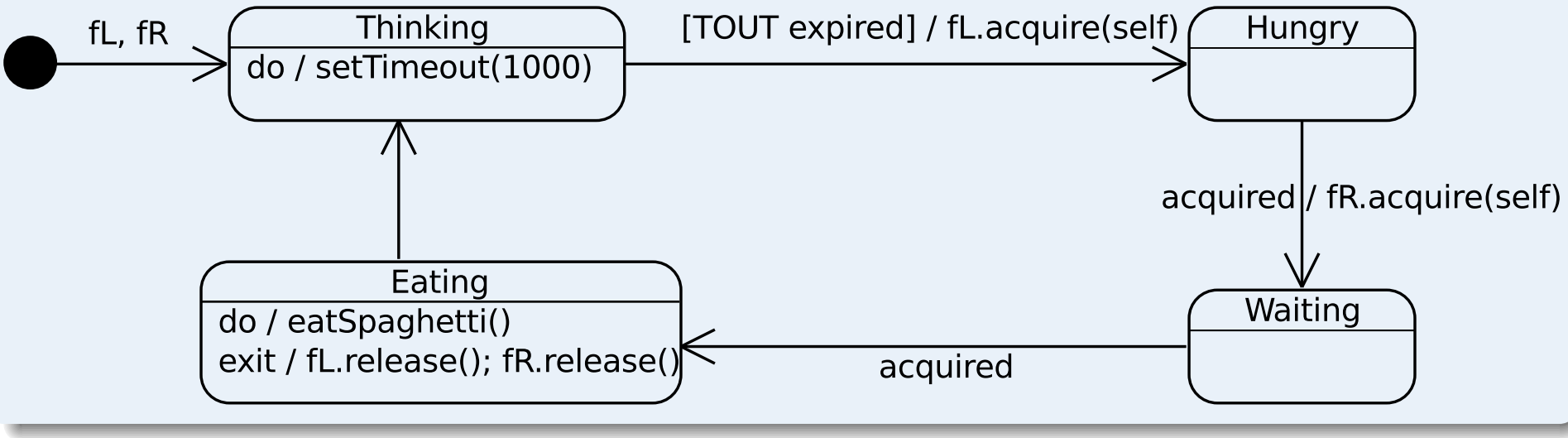
Fork



Please note: thinking time and fork grabbing order

A Transformation Approach: The Dining Philosophers (II)

Philosopher



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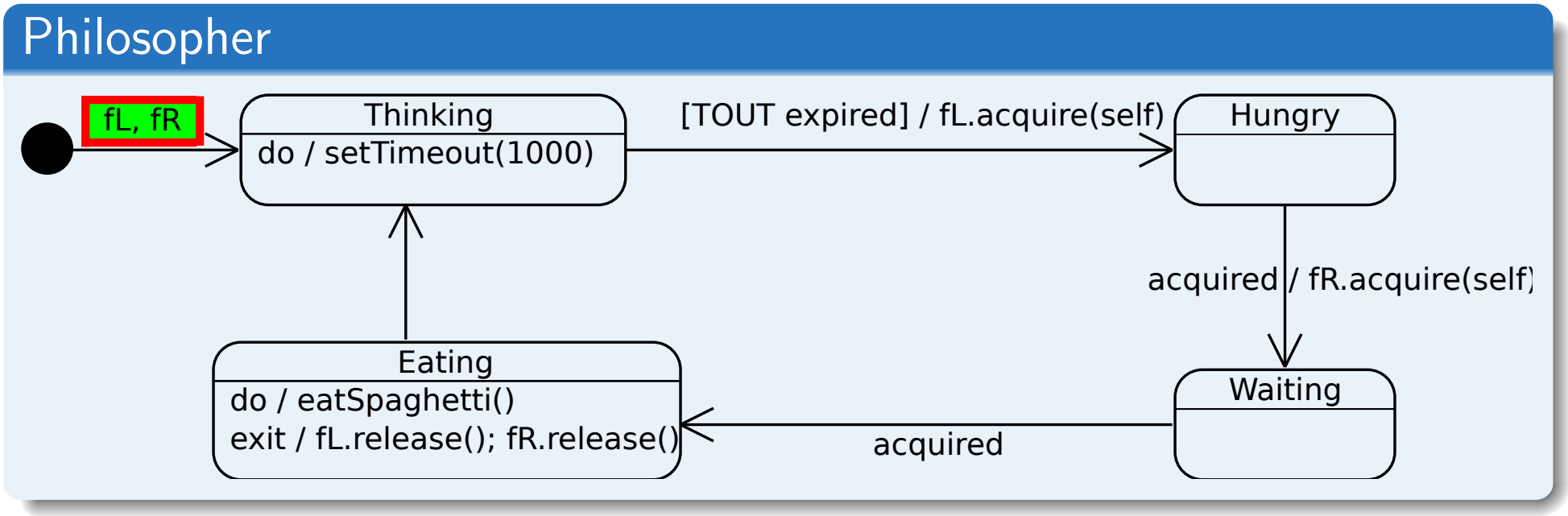

$$\mathcal{P} = \{fL, fR\}$$
$$\mathcal{E} = \{acquired\}$$

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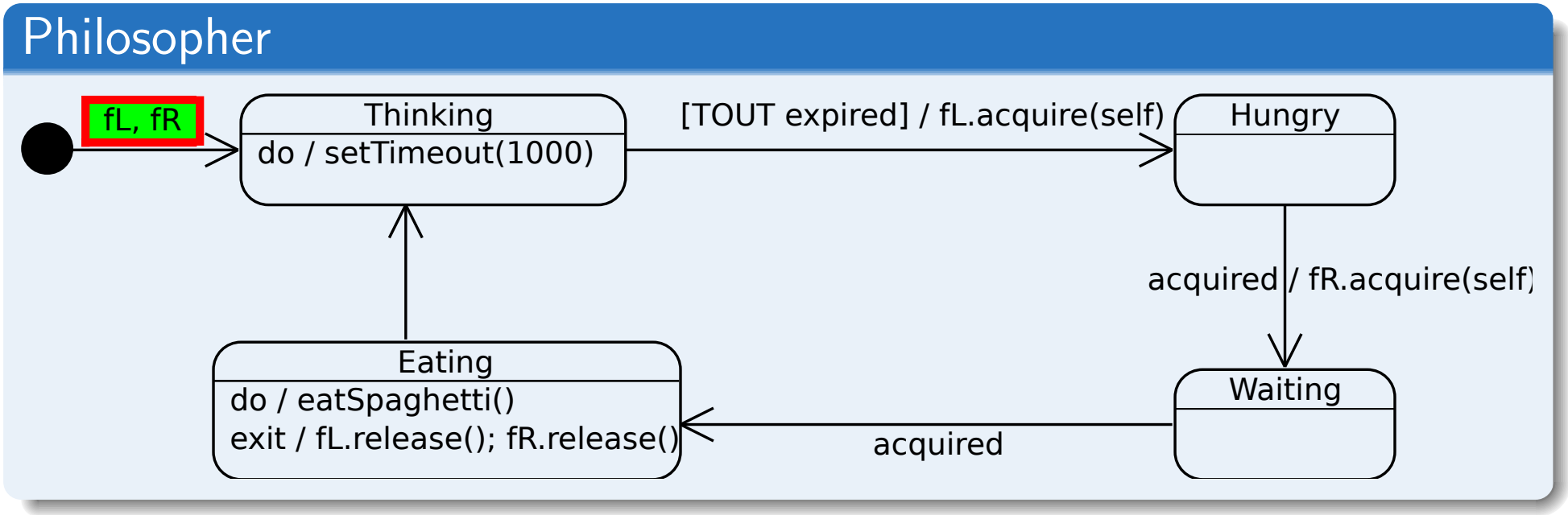

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-module(philosopher).
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start(FL, FR) ->
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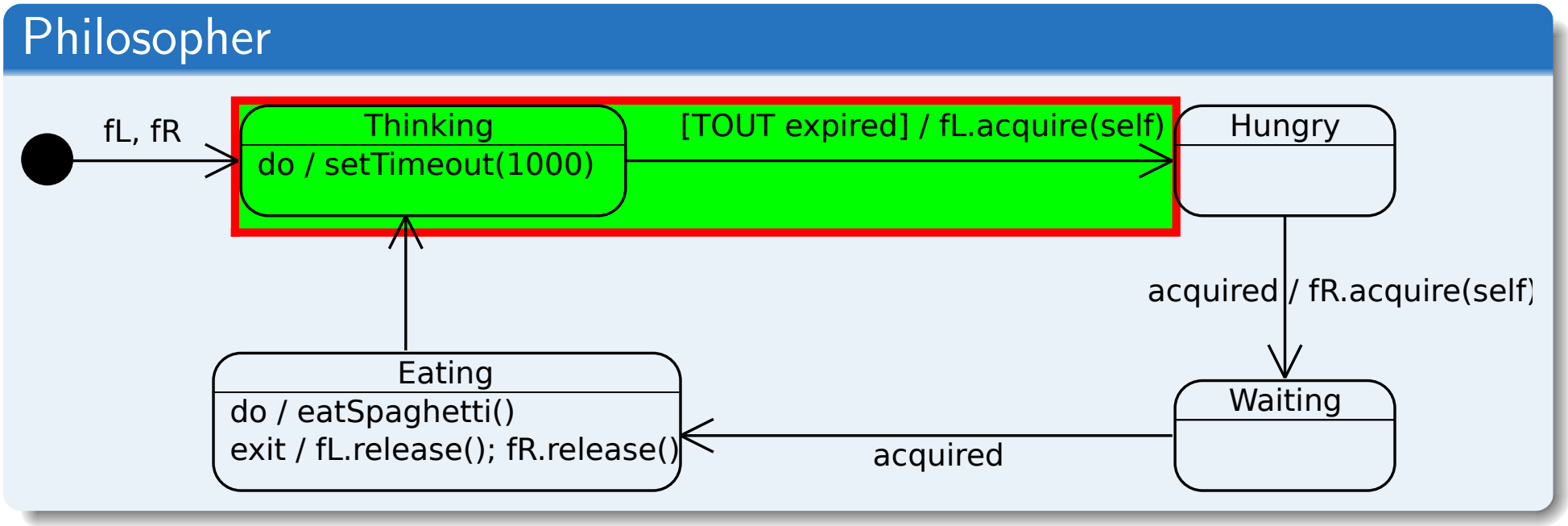

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```

```
thinking(FL, FR) ->
  receive
    X -> thinking(FL, FR)
    after 1000 ->
      FL!{acquire, self()}, hungry(FL, FR)
  end.
```

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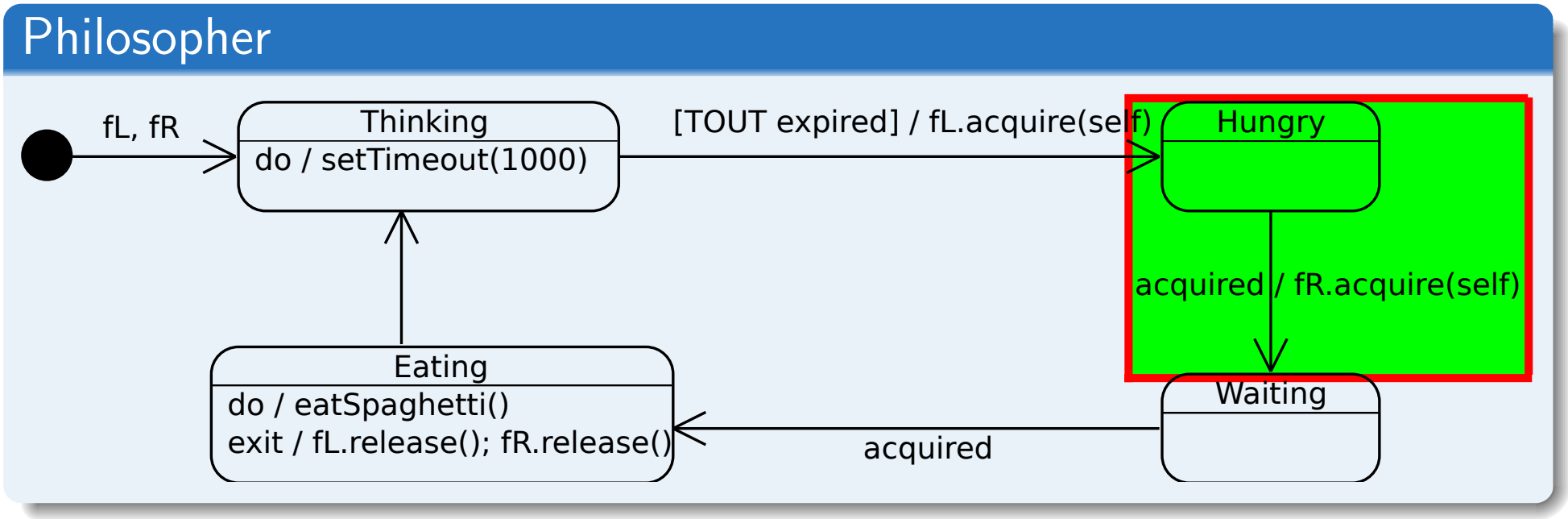

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thinking(FL, FR) -> ...
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```
hungry(FL, FR) ->
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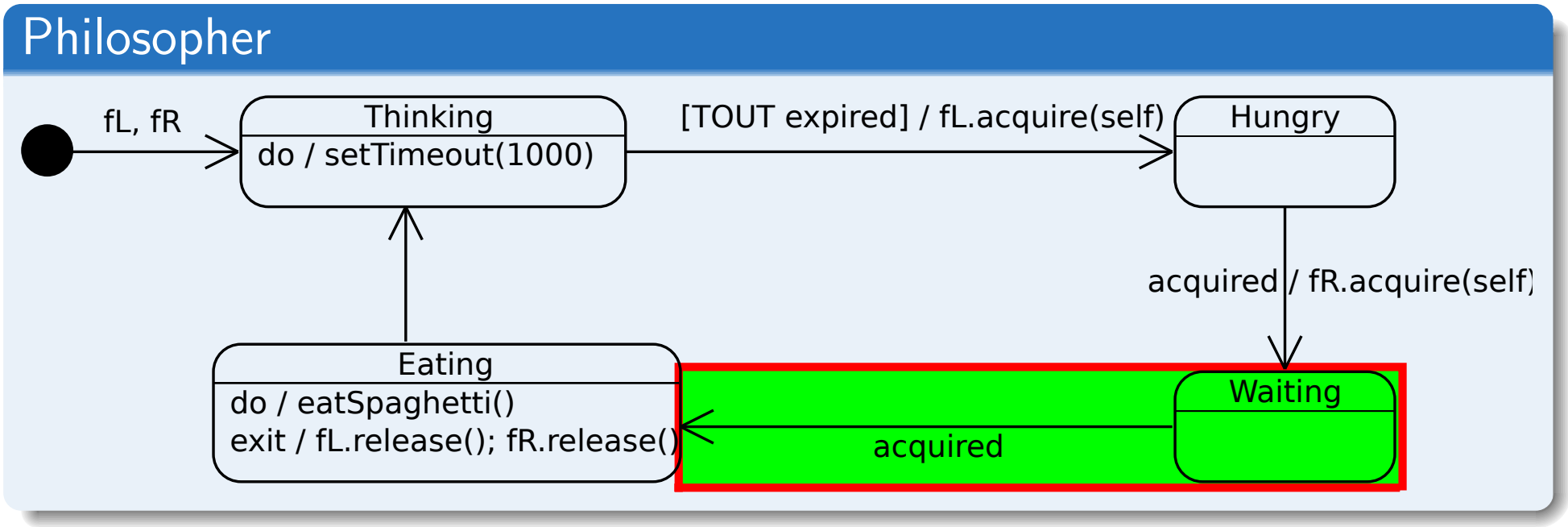

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waiting(FL, FR) ->
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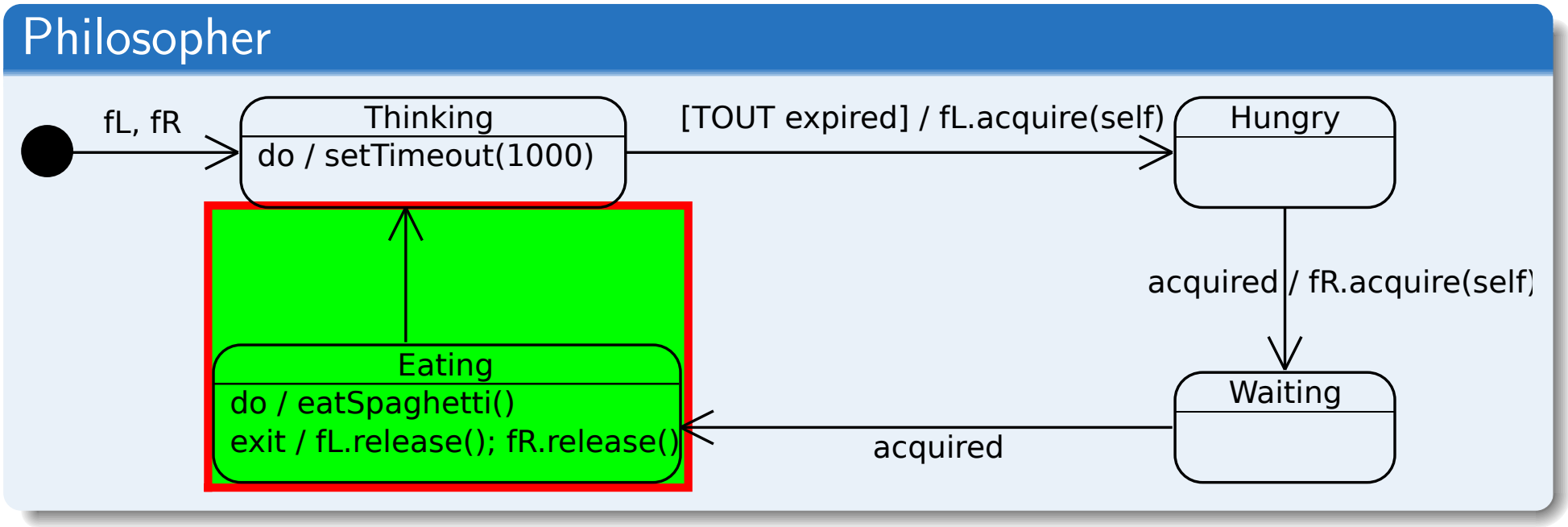

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thinking(FL, FR) -> ...
hungry(FL, FR) -> ...
waiting(FL, FR) -> ...
```

```
eating(FL, FR) ->
  eatSpaghetti(),
  FL!release, FR!release,
  thinking(FL, FR).
```

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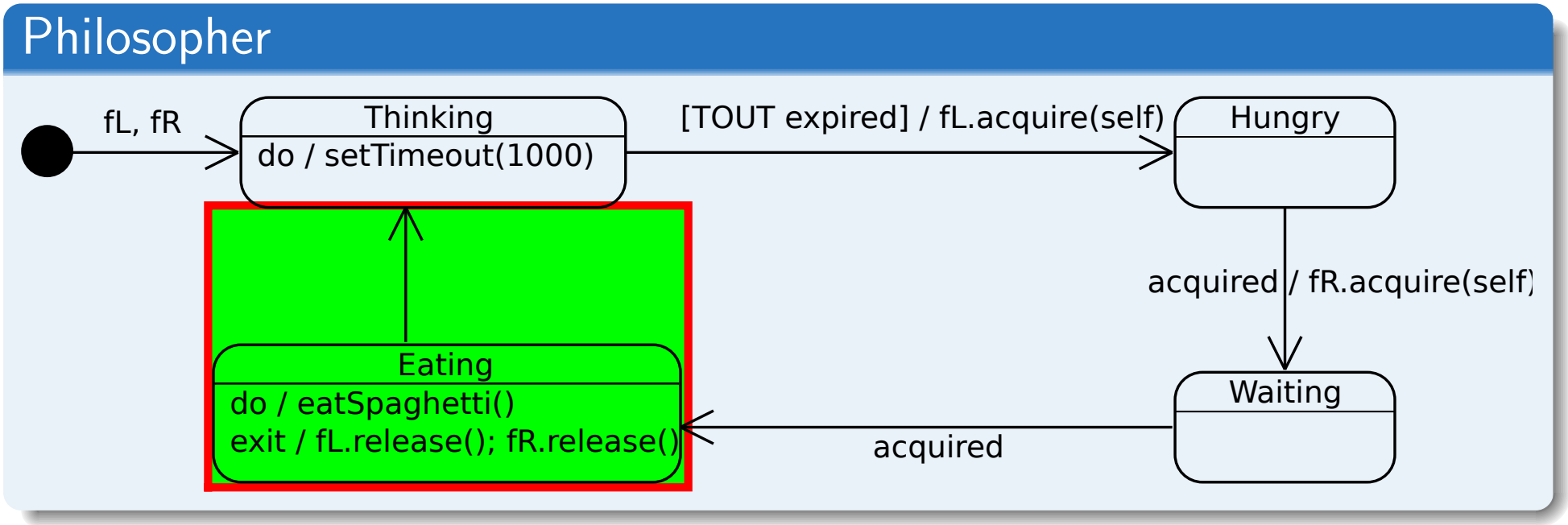

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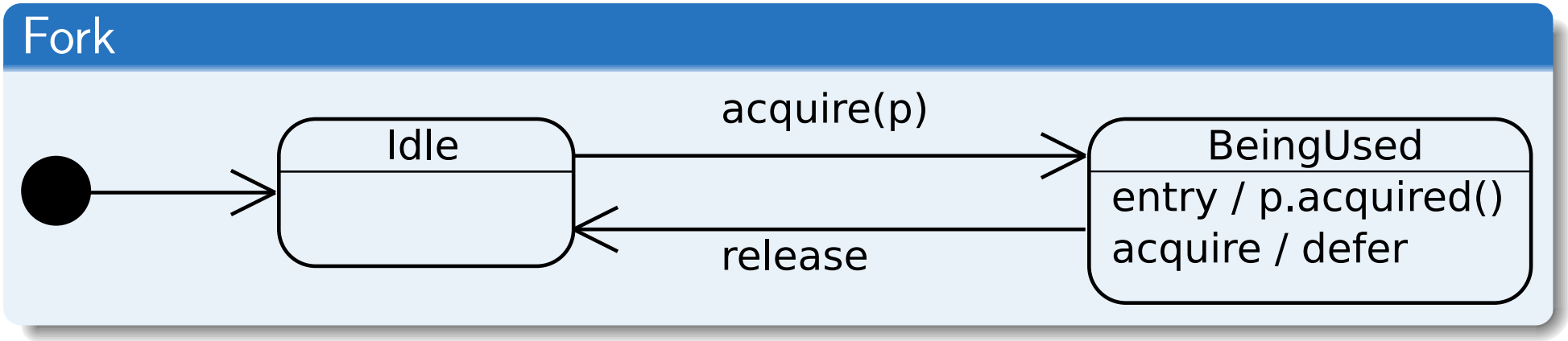
A UML instance is an Erlang process

A Transformation Approach: The Dining Philosophers (III)



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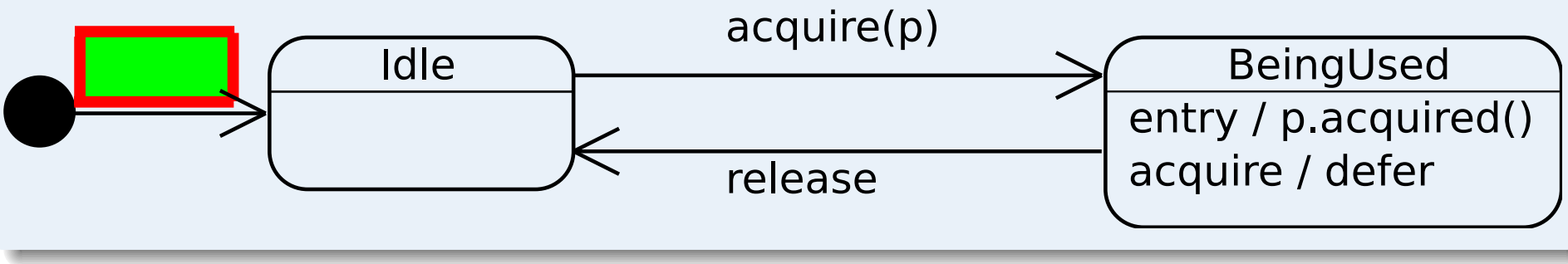

$$\mathcal{P} = \emptyset$$
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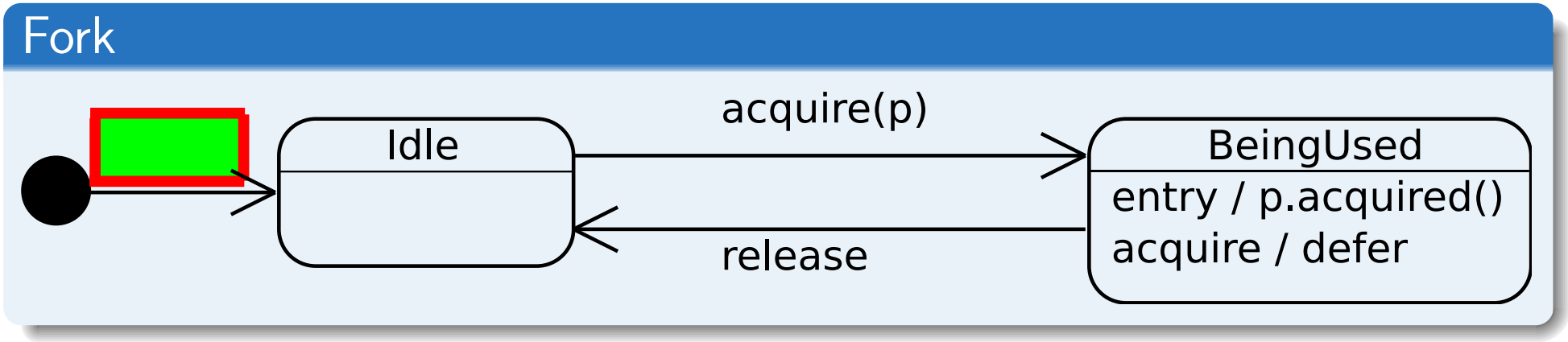

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-module(fork).
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start() ->
   spawn(fun() -> idle() end).
```

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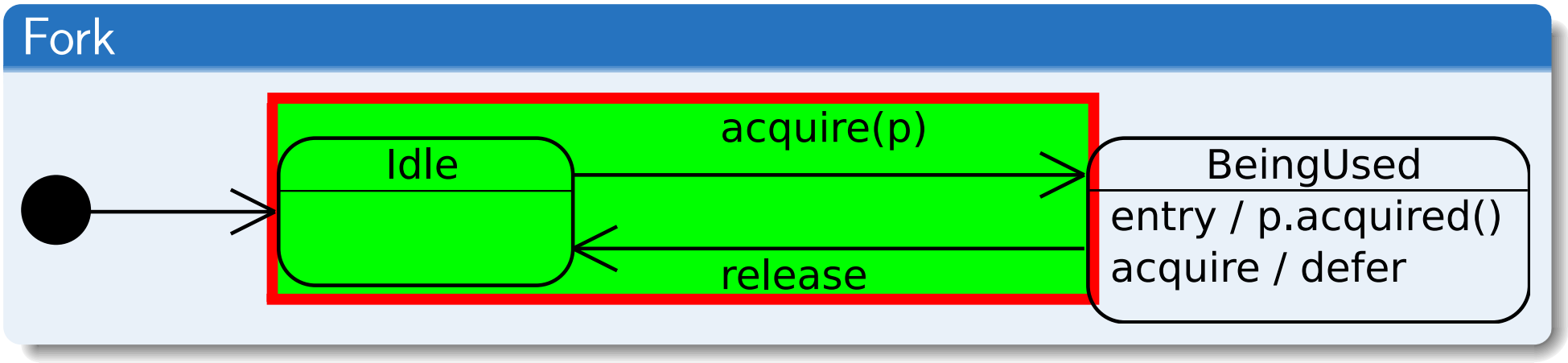

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-module(fork).
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```

```
idle() ->
  receive
    {acquire, P} -> beingUsed(P);
    X -> idle()
  end.
```

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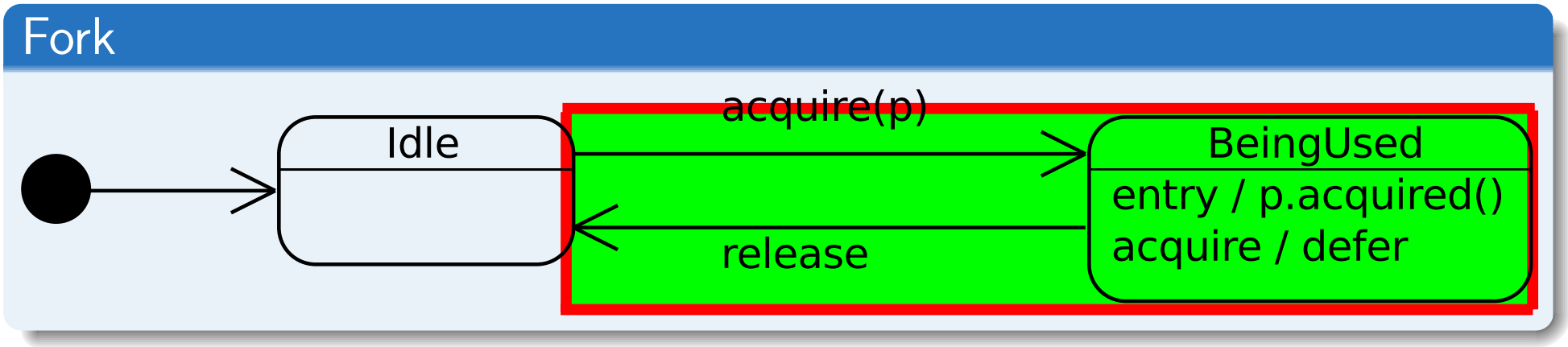

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A Transformation Approach: The Dining Philosophers (IV)

Wait! Explain me about `defer`...

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Assume that current state is `BeingUsed`, and `acquire` event is received

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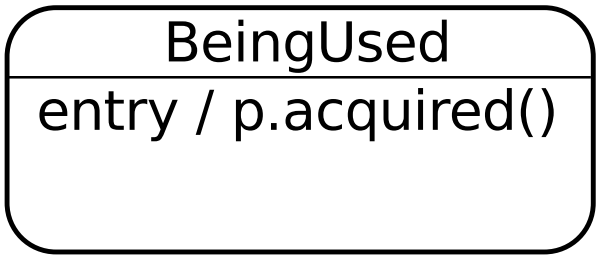

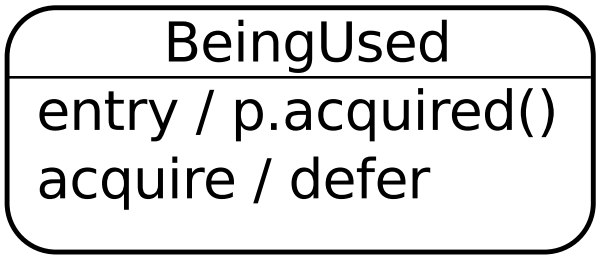

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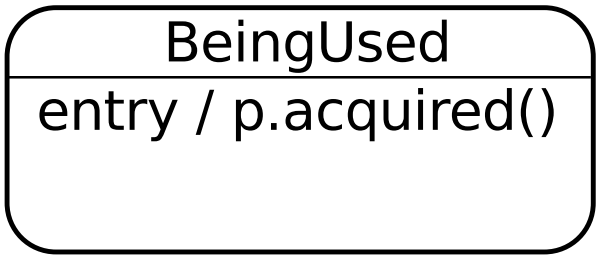

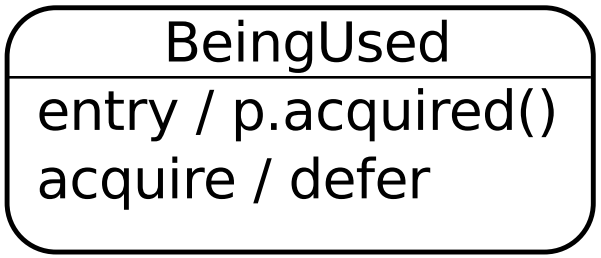

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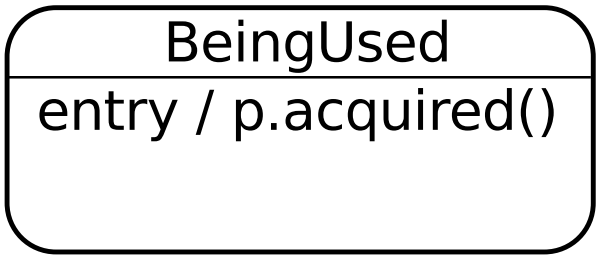

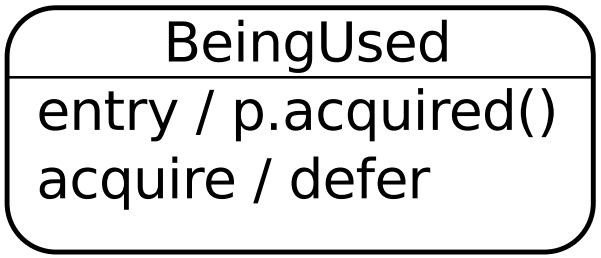

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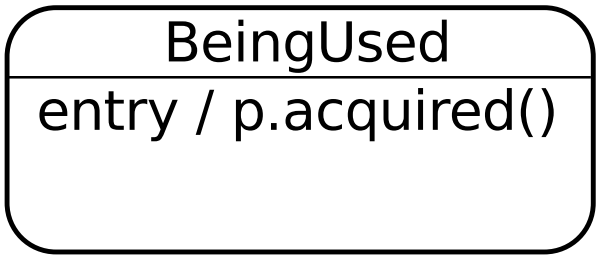

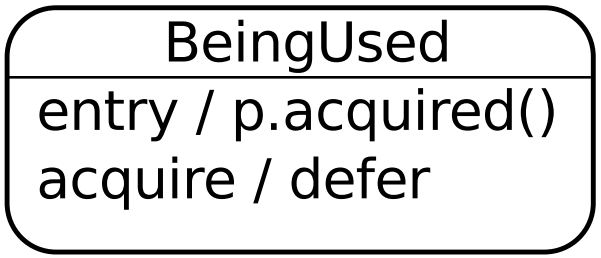

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Event has been discarded!

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Event is (eventually) handled

A Transformation Approach: The Dining Philosophers (IV)

Example system startup

```
run(N) ->
  Forks = lists:map (fun (_) -> fork:start() end, lists:
  lists:foreach
    (fun ({L,R}) -> philosopher:start(L, R) end, adjacen

adjacent([])      -> [];
adjacent([X|Xs]) -> lists:zip([X] ++ Xs, Xs ++ [X]).
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Using McErlang to verify correctness

```
> mce:start
    (#mce_opts{program=fun () -> dining:run(2) end,
               monitor=mce_mon_deadlock}).
...
*** Monitor failed
monitor error:
deadlock
```

A Transformation Approach: Algorithm (VI)

- Input data: UML-SM
- Output data: Erlang source code

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Algorithm steps (abstractedly)

1. Store parameters of initial transition ($\mathcal{P}$)
2. Create the Erlang header (`module`, `export`, `start`)
3. Create a set of triggered events of current state ($\mathcal{E}$)
4. Iterate for each state in the UML-SM
   1. Convert `entry`, `do` activities to message passing
   2. Special case: timeout activities
   3. Iterate in the output transitions
      1. Fill a `receive` Erlang skeleton properly
   4. Convert `exit` activity to message passing

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Related Work (I)

Automatic code generation

(multi-threaded behaviour and asynchronous communication)

Translator compiler

- C code from finite state machines with a synchronous semantics
  - PM-FORMS-03,AFLTY-ISORC-10
- Aynchronous semantics with a state table to reacts to events
  - NT-SEA-03,KNNZ-ICSE-00
- Design pattern forms implementing state machines
  - TKUY-ICRA-01
- Java thread per state-chart and Java objects to represent event queues
  - KM-TOOLS-02

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Interpreter to manage multi-threading and event queues $\rightarrow$ Erlang

Related Work (II)

Other model checkers

- Branching time model-checking using JACK
  - GLM-HASE-99
- Linear-time model checking using PROMELA
  - LMM-FAC-99
- UML class diagrams, UML-SMs and UML Communication diagrams verified using Maude LTL
  - CEC-IJSEA-12

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Conclusions and Future Work (I)

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- UML: standard as semi-formal modelling language
  - UML-SM: models system dynamics and its interaction
    - Modelling of concurrent and distributed systems
- Erlang: functional language
  - Good support for concurrency and distribution

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UML-SM $\rightarrow$ Erlang code

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$$\text{UML-SM} \rightarrow \text{Erlang code}$$

Contributions

- Reduce development time
  - By automatically generating Erlang skeleton code
- Enables validation of UML-SMs at an early development stage
  - Erlang-based model checking and testing techniques

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Contributions

- Reduce development time
  - By automatically generating Erlang skeleton code
- Enables validation of UML-SMs at an early development stage
  - Erlang-based model checking and testing techniques
- Also an alternative to Erlang code for "behaviour" pattern

Conclusions and Future Work (II)

Future Work

- Extend to additional UML-SM constructs
  - Preemptive UML-SM activities
  - Substates
  - Entry, exit or alternative potins
  - ...
- Apply to some real examples
- Tool support
  - Plugin for some UML CASE tool (e.g. Eclipse)

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