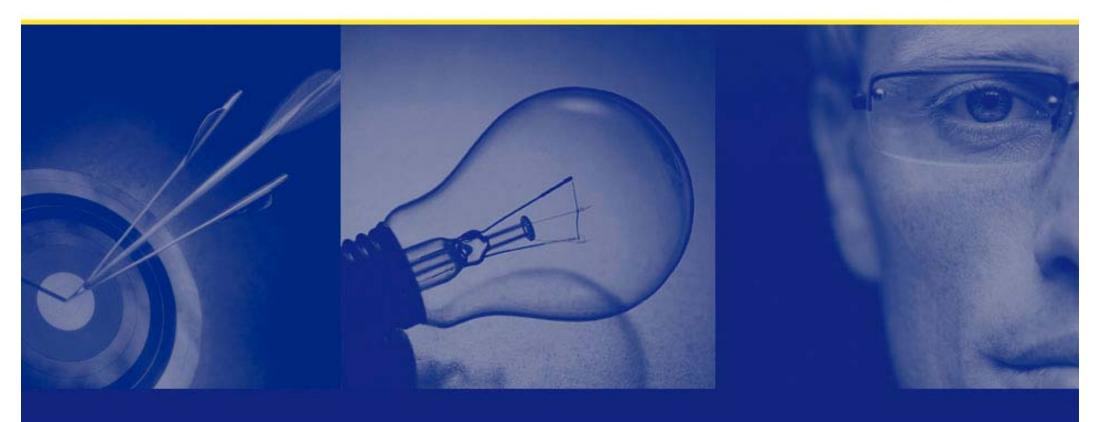
SQL Tuning via Toad





Tips for Optimizing SQL Performance

Bert Scalzo ...

Database Expert & Product Architect for Quest Software



Oracle Background:

- Worked with Oracle databases for over two decades (starting with version 4)
- Work history includes time at both "Oracle Education" and "Oracle Consulting"

Academic Background:

- Several Oracle Masters certifications
- BS, MS and PhD in Computer Science
- MBA (general business)
- Several insurance industry designations

Key Interests:

- Data Modeling
- Database Benchmarking
- Database Tuning & Optimization
- "Star Schema" Data Warehouses
- Oracle on Linux and specifically: RAC on Linux

Articles for:

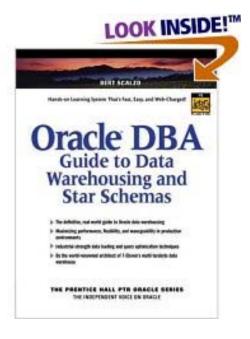
- Oracle's Technology Network (OTN)
- Oracle Magazine,
- Oracle Informant
- PC Week (eWeek)

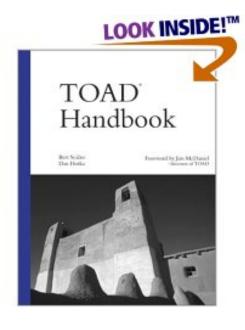
Articles for:

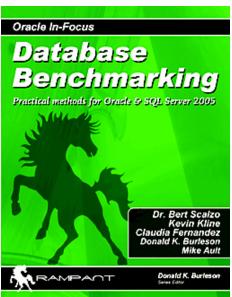
- Dell Power Solutions Magazine
- The Linux Journal
- www.linux.com
- www.orafaq.com

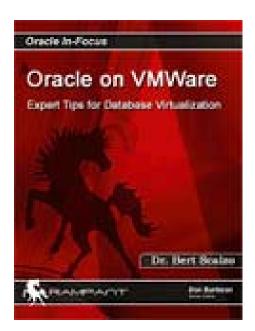


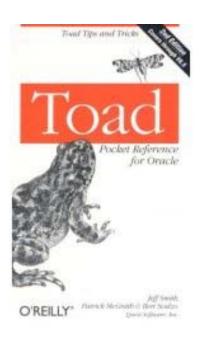
Books by Bert ...



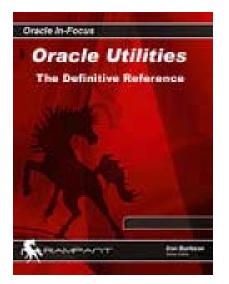








<u>Coming in 2009</u> ...



Also: <u>FREE</u> Toad e-Book for Toad 10...

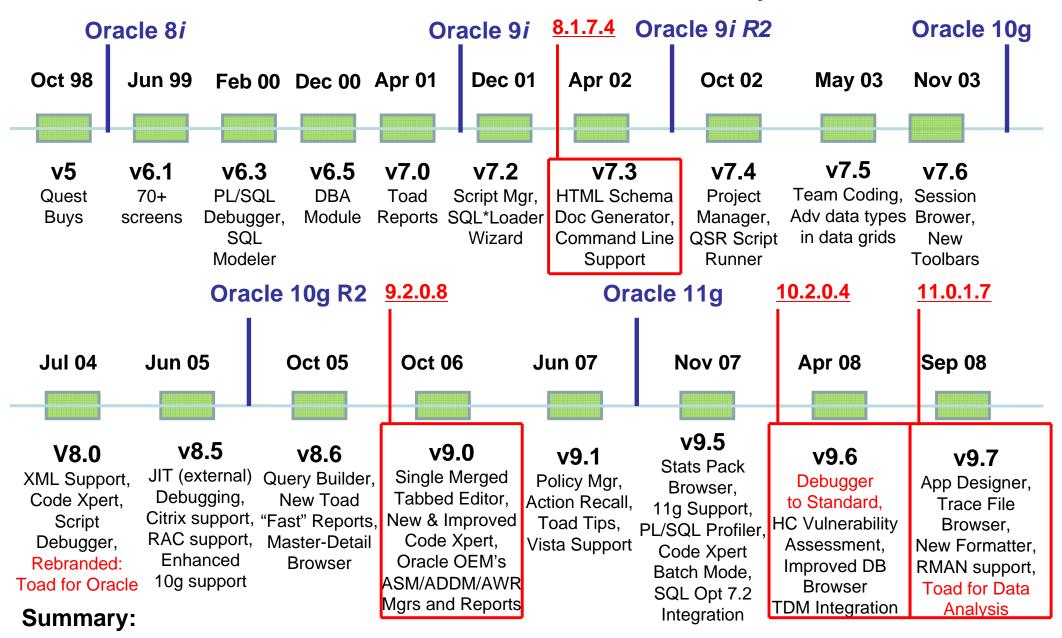


Topics ...

- Pre-Reqs
 - Correct Toad vs. Oracle Database Server version.
 - Correct Oracle SQL*Net Client networking version
 - SQL Tuning Approach much more than just explain plans and run times
- Explain Plans
 - Setup and effective use of the "Explain Plan"
 - Be careful, Explain Plan costs can sometimes not be the best way to pick the winner - sometimes (auto) trace is required to be 100% sure
 - Some guidelines on how to best or at least more easily read SQL explain plans - which is the general starting point for any SQL tuning attempt
- SQL Tuning Rules
 - Some Guidelines i.e. ("Golden Rules") just the tip of the iceberg
 - Efficient and fast selects & sub selects
 - Dealing with large tables
 - Parallel Hints
 Pinning SQL in Memory
 - Efficient SQL queries that use a lot of AND conditionals or sub-queries
 - How to avoid full-table scans
- Is There a Better (i.e. more productive) Way to Tune SQL
 - SQL Optimzier automate all the above (and much more)



Toad vs. Oracle Product Release History



Oracle 9i >= Toad 9.0 Oracle 10g >= Toad 9.6 Oracle 11g >= Toad 9.7 OUEST SOFTWARE

Oracle Client / Server Interoperability Support

(See Metalink Document 207303.1)

		Server Version									
Client Version	11.1.0	10.2.0	10.1.0	9.2.0	9.0.1	8.1.7	8.1.6	8.1.5	8.0.6	8.0.5	7.3.4
11.1.0	Yes	Yes #6	Yes #6	ES #5	No	No	No #3				
10.2.0	Yes #6	Yes	Yes	ES #5	No	Was	No #3				
10.1.0 <u>(#4)</u>	Yes #6	Yes	Yes	ES	Was	Was #2	No #3				
9.2.0	ES #5	ES #5	ES	ES	Was	Was	No	No	Was	No	No #1
9.0.1	No	No	Was	Was	Was	Was	Was	No	Was	No	Was
8.1.7	No	Was	Was	Was	Was	Was	Was	Was	Was	Was	Was
8.1.6	No	No	No	No	Was	Was	Was	Was	Was	Was	Was
8.1.5	No	No	No	No	No	Was	Was	Was	Was	Was	Was
8.0.6	No	No	No	Was	Was	Was	Was	Was	Was	Was	Was
8.0.5	No	No	No	No	No	Was	Was	Was	Was	Was	Was
7.3.4	No	No	No	Was	Was	Was	Was	Was	Was	Was	Was

Key:

Yes	Supported
ES	Supported but fixes only possible for customers with Extended Support.
Was	Was a supported combination but one of the releases is no longer covered by any of Premier Support, Primary Error Correct support, Extended Support nor Extended Maintenance Support so fixes are no longer possible.
No	Has never been Supported

Toad may work with older client talking to newer databases - but there might be data type issues ⊗



Seven Steps for SQL Tuning Success

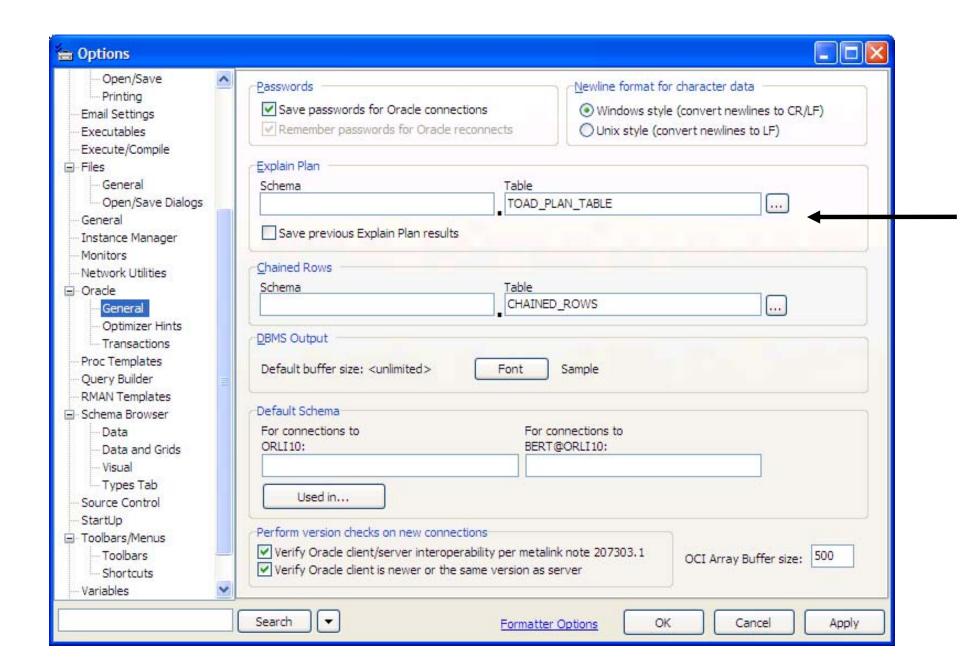
- 1. Always start by knowing (i.e. being able to say in English) what the query does
- 2. For queries involving more than 2 tables, a data model can be a handy road map
- 3. Explain plan costs alone may well lead you astray sometimes the costs can lie
- 4. Sometimes equal execution times don't necessarily equate to equivalent solutions
- 5. You should always include (auto) trace information to divine among all the above
- 6. Sole reliance on automatic SQL optimization and tuning tools can be suboptimal
- 7. You must add human intuition and insight to the optimization process for success



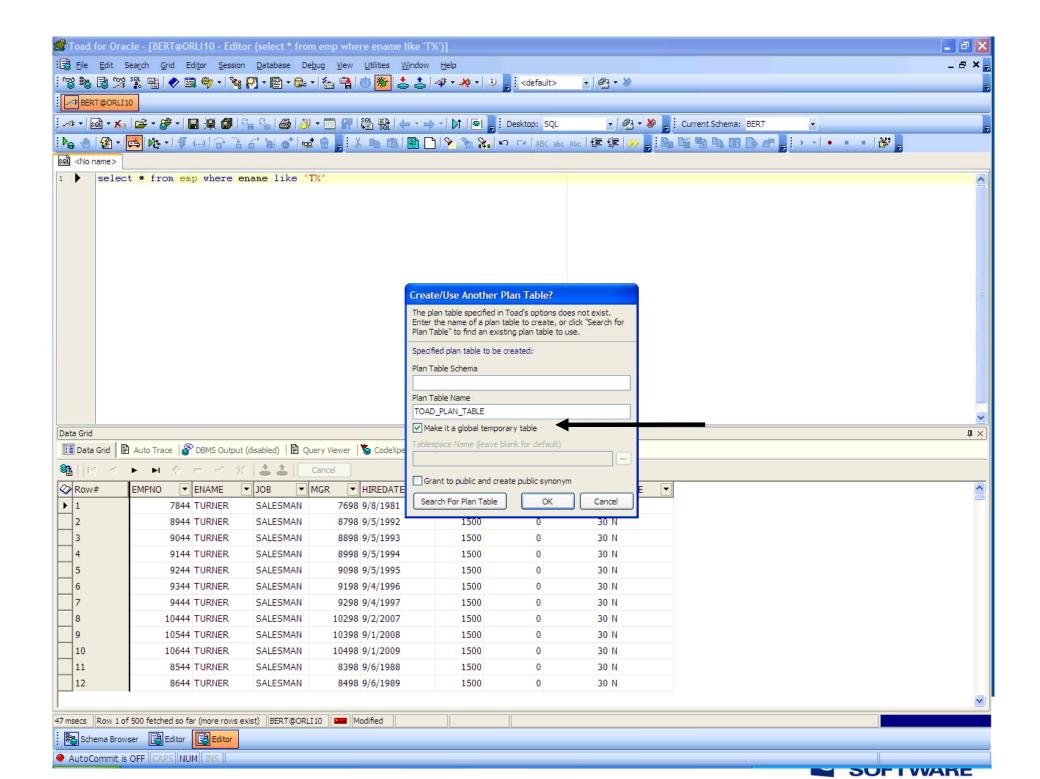
Explain Plans

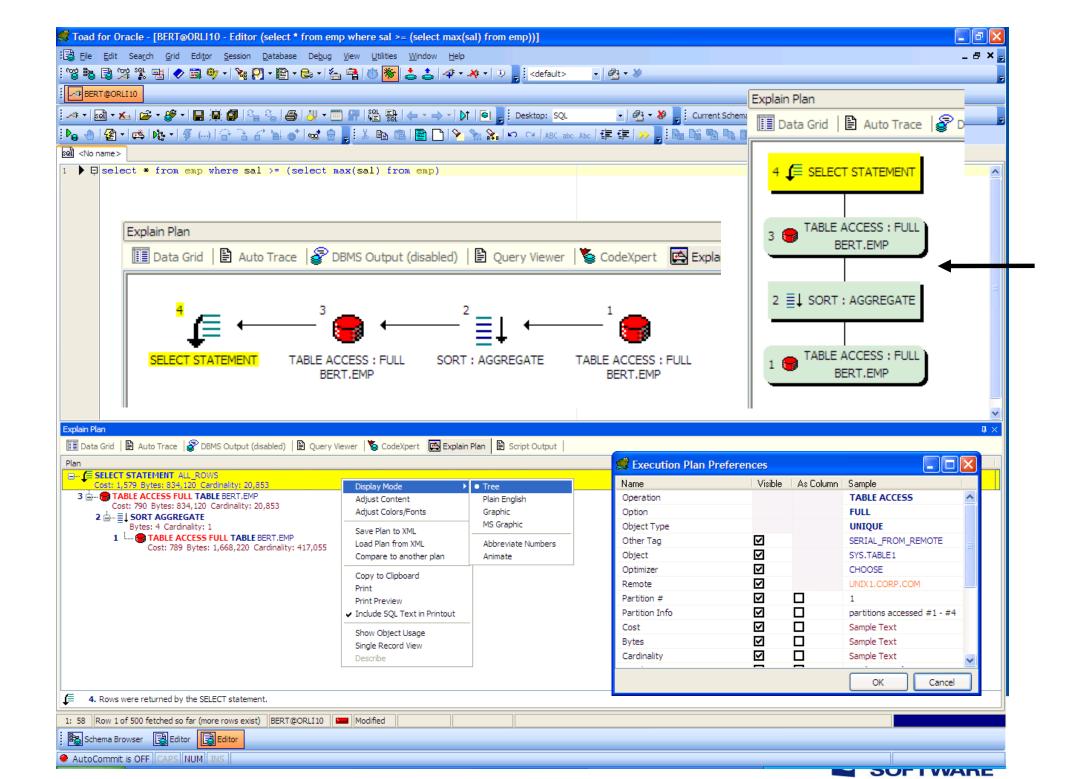
- Explain Plans are the standard Oracle mechanism to peek into the possible "internal algorithm" the database engine might execute for the query (think of it as sort of like program pseudo-code)
- Explain Plans generally require an Oracle "plan table" to hold the explain plan intermediate results
 - Three Options here:
 - Central "plan table" for all users to share managed by DBA
 - "Plan table" per schema but be careful if users all login the same
 - "Plan table" per session -
- When doing explain plans manually
 - Method #1
 - EXPLAIN FOR SELECT * FROM emp;
 - SELECT ... FROM plan_table WHERE ... (fairly complex SQL)
 - Method #2
 - EXPLAIN FOR SELECT * FROM emp;
 - SELECT * FROM table(DBMS_XPLAN.DISPLAY(PLAN_TABLE));

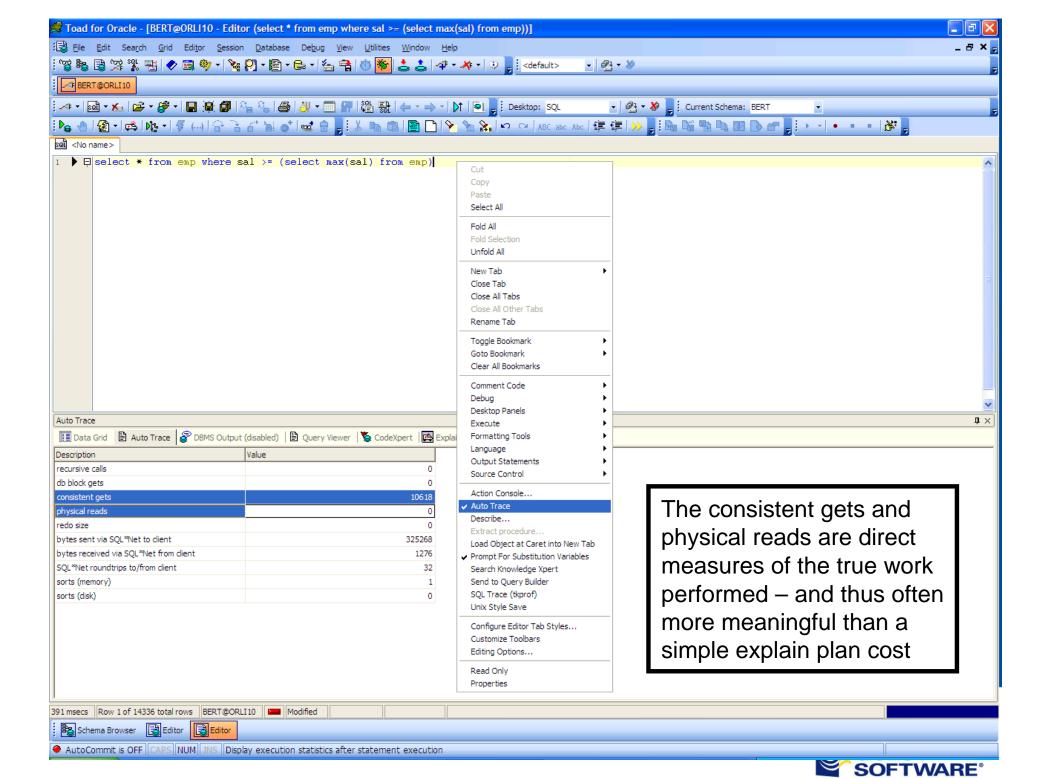


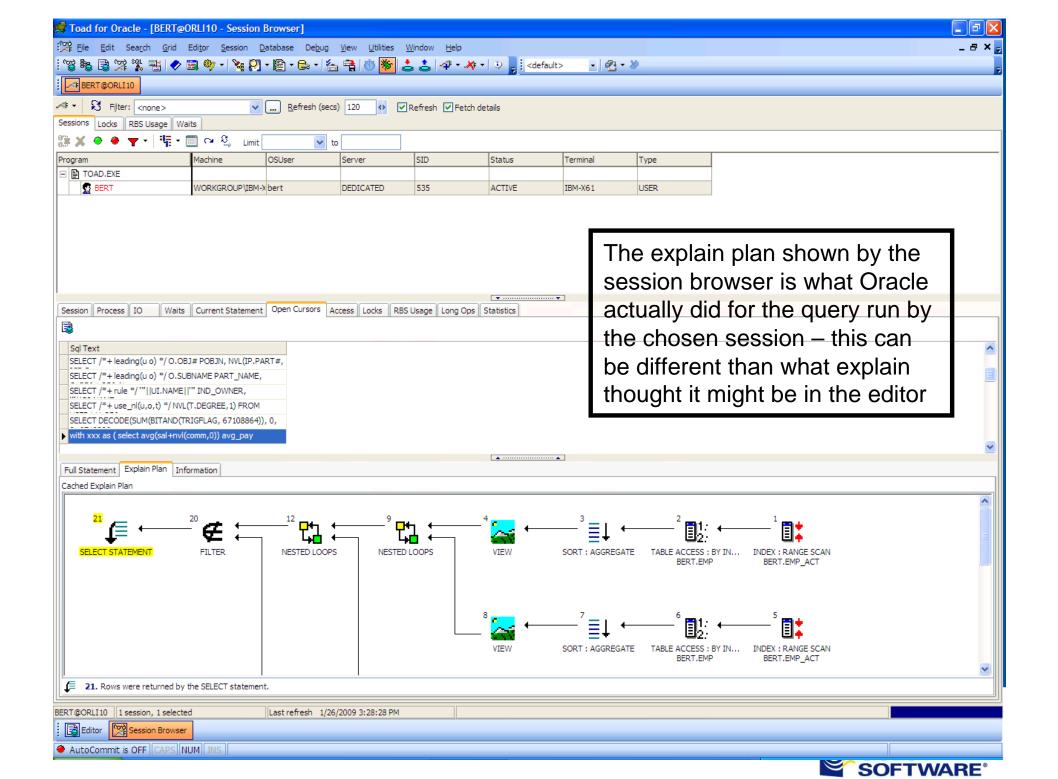












Rule #1: Watch Indexed WHERE Conditions

Assume <u>address index</u> has columns (city, state)	
•non-leading index column references may not use indexes	
•where state = 'TX'	[Depends Oracle on Version]
•where city = 'DALLAS'	[Index Used]
•where state = 'TX' and city = 'DALLAS'	[Index Used]
•NOT, != and <> disable index use	
•where state not in ('TX', 'FL', 'OH')	[Index Not used]
•where state != 'TX'	[Index Not used]
•NULL value references almost never use indexes (one exception - bitma	ps)
•where state IS NULL	[Index Not used]
•where state IS NOT NULL	[Index Not used]
•expression references can never use indexes	
•where substr(city,1,3) = 'DAL'	[Index Not used]
•where city like 'DAL%'	[Index Used]
•where city state = 'DALLASTX'	[Index Not used]
•where city = 'DALLAS' and state = 'TX'	[Index Used]
•where salary * $12 >= 24000$	[Index Not used]
•where salary >= 2000	[Index Used]



Rule #2: Watch Non-Indexed WHERE Conditions

- •Oracle evaluates Non-Indexed conditions linked by "AND" bottom up
 - •Bad: select * from address where

```
areacode = 972 and
```

type_nr = (select seq_nr from code_table where type = 'HOME')

•Good: select * from address where

type_nr = (select seq_nr from code_table where type = 'HOME') and areacode = 972

- •Oracle evaluates Non-Indexed conditions linked by "OR" top down
 - •Bad: select * from address where

type_nr = (select seq_nr from code_table where type = 'HOME') **or** areacode = 972

•Good: select * from address where

areacode = 972 or

type_nr = (select seq_nr from code_table where type = 'HOME')



Rule #3: Order Table in the FROM Clause (pre-10g)

- •important under rule based optimizer, and won't hurt under cost based optimizer
- •order FROM clauses in descending order of table sizes based upon row counts
- •for example
 - •select * from larger table, smaller table
 - •select * from larger table, smaller table, smallest table
 - •select * from larger table, smaller table, associative table

Note – rule based optimizer only (pre-10g)



Rule #4: Consider IN or UNION in place of OR

```
•if columns are not indexed, stick with OR
•if columns are indexed, use IN or UNION in place of OR
            •IN example
                  •Bad: select * from address where
                                     state = TX' or
                                     state = 'FL' or
                                     state = 'OH'
                  •Good: select * from address where
                                     state in ('TX','FL','OH')
            •UNION example
                  •Bad: select * from address where
                                     state = 'TX' or
                                     areacode = 972
                  •Good: select * from address where
                                     state = TX
                          union
                          select * from address where
                                     areacode = 972
```



Rule #5: Weigh JOIN versus EXISTS Sub-Query

- •use table JOIN instead of EXISTS sub-query
 - •when the percentage of rows returned from the outer sub-query is high

```
select e.name, e.phone, e.mailstop
from employee e, department d
where e.deptno = d.deptno
and d.status = 'ACTIVE'
```

- •use EXISTS sub-query instead of table JOIN
 - •when the percentage of rows returned from the outer sub-query is low

```
select e.name, e.phone, e.mailstop
from employee e
where e.deptno in (select d.deptno
from department d
where d.status != 'ACTIVE')
```



Rule #6: Consider EXISTS in place of DISTINCT

•avoid joins that use DISTINCT, use EXISTS sub-query instead

•**Bad:** select **distinct** deptno, deptname from emp, dept where emp.deptno = dept.deptno

•Good: select deptno, deptname from dept where

exists (select 'X' from emp where

emp.deptno = dept.deptno)

Note – only has to find one match



Rule #7: Consider NOT EXISTS in place of NOT IN

•avoid sub-queries that use NOT IN, use NOT EXISTS instead

•Bad: select * from emp where

deptno **not in** (select deptno from dept where

deptstatus = 'A')

•Good: select * from emp where

not exists (select 'X' from dept where

deptstatus = 'A' and

dept.deptno = emp.deptno)

Note – only has to find one non-match



Rule #8: Ordering Via the WHERE Clause

•a dummy WHERE clause referencing an indexed column will

•retrieve all records in ascending order (descending for 8i descending index)

•not perform a costly sort operation

•Bad: select * from address order by city

•Good: select * from address where city > "



Rule #9: Use PL/SQL to reduce network traffic

```
•Utilize PL/SQL to group related SQL commands and thereby reduce network traffic
      •Bad:
            select city_name, state_code
               into:v city,:v sate
               from zip_codes where zip_code = '75022';
            insert into customer ('Bert Scalzo', '75022', :v city, v state);
      •Good:
            begin
               select city_name, state_code
                 into:v_city,:v_sate
                 from zip_codes where zip_code = '75022';
               insert into customer ('Bert Scalzo', '75022', :v_city, v_state);
            end;
```



Rule #10: Partition Large Tables and Indexes

- Partition Elimination
- •Partition-Wise Join (requires Parallel too)
- •NOTE: Do not expect that merely partitioning will solve some major performance problem, it should merely make an incremental improvement to a non-partitioned explain plan. Read that as partitioning can make an already good explain plan even better.



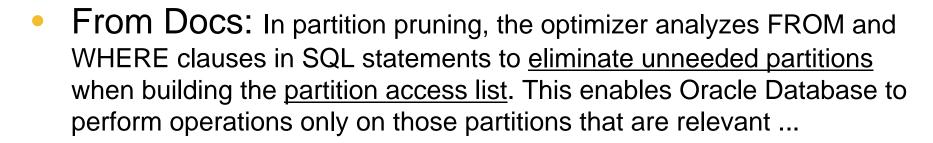
Partitioning Benefits: Opinion (Mine)

•	Performance	20%	
•	Capacity Management	20%	Tartition
•	Availability	20%	Why to Partition
•	Manageability	40%	\//b>/ 40

- Don't over-sell/over-expect the performance aspect
- Need to experiment for best approach for a database
- Better to take longer at the start to get right, because very often it's far too expensive to change afterwards



Partition Pruning (Restriction Based) 🖈



- "Divide and Conquer" for performance
 - Sometimes can yield order of magnitude improvement
 - But once again, best not to oversell and/or over-expect
- Some Potential Issues to be aware of:
 - SQL*Plus Auto-Trace can sometimes miss partition pruning
 - "Old Style" Explain Plans via simple SELECT has issues too
 - Best to always use <u>DBMS_XPLAN</u> and/or <u>SQL_TRACE</u>



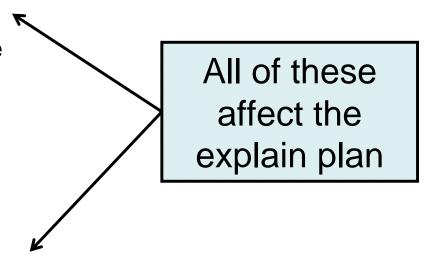
Partition-Wise Join (Multi-Object Based)



From Docs: Partition-wise joins reduce query response time by minimizing the amount of data exchanged among parallel execution servers when joins execute in parallel. This significantly reduces response time & improves the use of both CPU & memory resources.

Different Flavors:

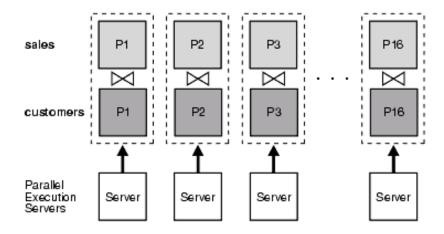
- Full Single to Single
- Full Composite to Single
- Full Composite to Composite
- Partial Single
- Partial Composite
- **Indexing Strategy Counts**
 - Local Prefixed/Non-Prefixed
 - Global



Picture Worth 1000 Words (from Docs)

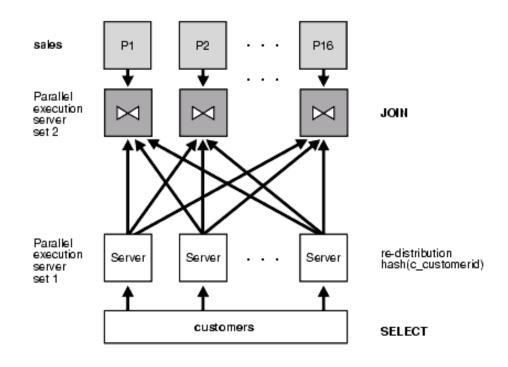
Simple Mantra: Subdivide the work into equally paired chunks, then perform all that work using many parallel processes

Figure 4-1 Parallel Execution of a Full Partition-wise Join



Make sure not to over-allocate CPU's – remember there will also be concurrent workload

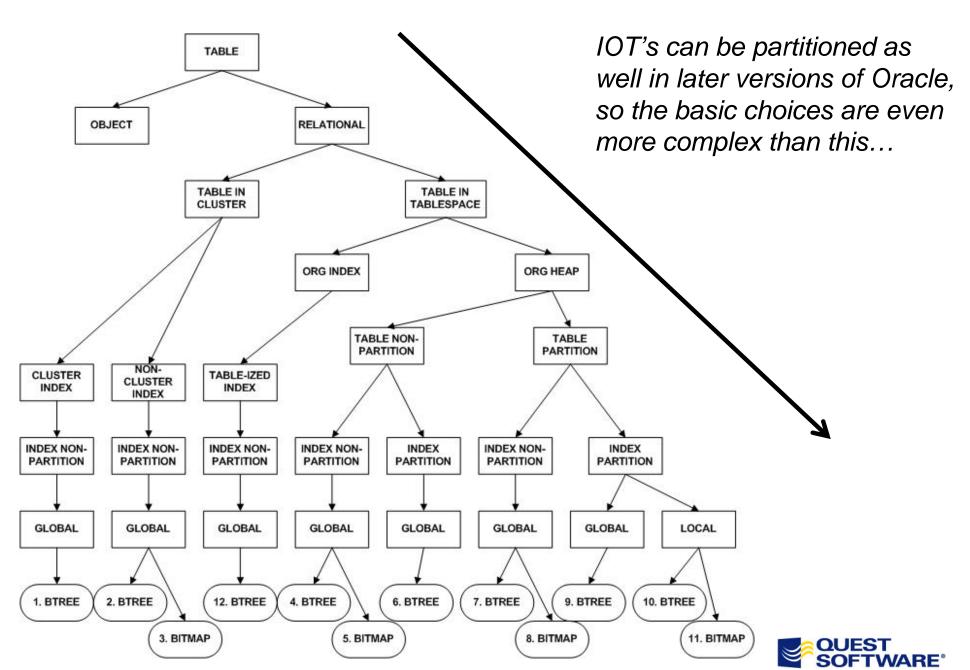
Figure 4-3 Partial Partition-Wise Join



Partitioning History (from Oracle 11G training+)

Oracle 5	Before Tablespac	es – we had pa	artitions ☺ ←				
Oracle 7	<u> </u>	Partition Views – really more of a cheat ⊗ ←					
	Core functionality	Core functionality Performance Manageability					
Oracle8	Range partitioning Global range indexes	"Static" partition pruning	Basic maintenance operations: add, drop, exchange				
Oracle8i	Hash and composite range-hash partitioning	Partition-wise joins "Dynamic" pruning	Merge operation				
Oracle9i	List partitioning		Global index maintenance				
Oracle9i R2	Composite range-list partitioning	Fast partition split					
Oracle10g	Global hash indexes		Local Index maintenance				
Oracle10g R2	1M partitions per table	"Multi-dimensional" pruning	Fast drop table				
Oracle Database 11g	More composite choices REF Partitioning Virtual Column		Interval Partitioning Partition Advisor				

Partitioning Options – Part 1



Partitioning Options – Part 2

Prior to 11G: Oracle White Paper: 2007 Partitioning in Oracle Database 11g

Partitioning Strategy	Data Distribution	Sample Business Case
Range Partitioning	Based on consecutive ranges of values.	Orders table range partitioned by order_date
List Partitioning	Based on unordered lists of values.	Orders table list partitioned by country
Hash Partitioning	Based on a hash algorithm.	Orders table hash partitioned by customer_id
Composite Partitioning Range-Range Range-List Range-Hash List-List List-Range List-Hash	Based on a combination of two of the above-mentioned basic techniques of Range, List, Hash, and Interval Partitioning	 Orders table is range partitioned by order_date and sub-partitioned by hash on customer_id Orders table is range partitioned by order_date and sub-partitioned by range on shipment_date



Partitioning Options – Part 3

Post 11G: Oracle White Paper: 2007 Partitioning in Oracle Database 11g

Partitioning Extension	Partitioning Key	Sample Business Case	
Interval Partitioning • Interval • Interval-Range • Interval-List • Interval-Hash	An extension to Range Partition. Defined by an interval, providing equi-width ranges. With the exception of the first partition all partitions are automatically created on- demand when matching data arrives.	Orders table partitioned by order_date with a predefined daily interval, starting with '01-Jan-2007'	
REF Partitioning	Partitioning for a child table is inherited from the parent table through a primary key – foreign key relationship. The partitioning keys are not stored in actual columns in the child table.	(Parent) Orders table range partitioned by order_date and inherits the partitioning technique to (child) order lines table. Column order_date is only present in the parent orders table.	← Very exciting
Virtual column based Partitioning	Defined by one of the above- mentioned partition techniques and the partitioning key is based on a virtual column. Virtual columns are not stored on disk and only exist as metadata.	Orders table has a virtual column that derives the sales region based on the first three digits of the customer account number. The orders table is then list partitioned by sales region.	new options ←──────────────────────────────────

Operation	Name	Rows	Bytes	Cost	Pstart Ps	stop
SELECT STATEMENT		1	154	34	 	
SORT GROUP BY		1	154	34	İ	
HASH JOIN		1	154	29		
TABLE ACCESS BY INDEX ROWID	DW_ORDER	1	95	17		
BITMAP CONVERSION TO ROWIDS						
BITMAP AND						
BITMAP MERGE						
BITMAP KEY ITERATION						
TABLE ACCESS BY INDEX ROWID	DW_PERIOD	1	51	2		
BITMAP CONVERSION TO ROWIDS						
BITMAP AND						
BITMAP INDEX SINGLE VALUE	DW_PERIOD_B03					
BITMAP INDEX SINGLE VALUE	DW_PERIOD_B12					
BITMAP INDEX RANGE SCAN	DW_ORDER_B1					
BITMAP MERGE						
BITMAP KEY ITERATION						
TABLE ACCESS BY INDEX ROWID	DW_LOCATION	1	46	2		
BITMAP CONVERSION TO ROWIDS						
BITMAP AND						
BITMAP INDEX SINGLE VALUE	DW_LOCATION_B03					
BITMAP INDEX SINGLE VALUE	DW_LOCATION_B41					
BITMAP INDEX RANGE SCAN	DW_ORDER_B2					
BITMAP MERGE						
BITMAP KEY ITERATION						
TABLE ACCESS BY INDEX ROWID	DW_PRODUCT	17	1K	10		
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HASH JOIN	ĺ	1	154	30		
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TABLE ACCESS BY LOCAL INDEX ROWID	DW_ORDER_PART	0	20	18	1	10
BITMAP CONVERSION TO ROWIDS						
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BITMAP KEY ITERATION						
SORT BUFFER						
TABLE ACCESS BY INDEX ROWID	DW_PERIOD	1	51	2		
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BITMAP INDEX SINGLE VALU	:					
BITMAP INDEX RANGE SCAN	DW_ORDER_PART_B1				1	10
BITMAP MERGE						
BITMAP KEY ITERATION						
SORT BUFFER		1				
TABLE ACCESS BY INDEX ROWID	1	1	46	2		
BITMAP CONVERSION TO ROWID						
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Rule #11: Serial Explain Plans, then Parallel (maybe)

- •Parallel Full Table Scan
- Parallel Index Scan
- •Parallel Fast Full Scan (FFS Index Scan)
- •NOTE: Do not expect that merely parallelizing will solve some major performance problem, it should merely make an incremental improvement to a non-paralell (i.e. serial) explain plan. Read that as parallel can make an already good explain plan even better.



- •Parallel processing is controlled as follows:
 - •Query has /*+parallel*/ hint
 - •Some shops do NOT favor hints
 - •What if database version changes
 - •What happens if statistics change
 - •Other questionable future scenarios
 - •Cannot add hints to pre-canned applications
 - •Object (table or index) has parallel degree
 - •Database instance parameter for parallel
- •For RAC, parallel can also span the RAC nodes too



Operation	Name	Rows	Bytes	Cost	Pstart Ps	stop
SELECT STATEMENT		1	154	34	 	
SORT GROUP BY		1	154	34	İ	
HASH JOIN		1	154	29		
TABLE ACCESS BY INDEX ROWID	DW_ORDER	1	95	17		
BITMAP CONVERSION TO ROWIDS						
BITMAP AND						
BITMAP MERGE						
BITMAP KEY ITERATION						
TABLE ACCESS BY INDEX ROWID	DW_PERIOD	1	51	2		
BITMAP CONVERSION TO ROWIDS						
BITMAP AND						
BITMAP INDEX SINGLE VALUE	DW_PERIOD_B03					
BITMAP INDEX SINGLE VALUE	DW_PERIOD_B12					
BITMAP INDEX RANGE SCAN	DW_ORDER_B1					
BITMAP MERGE						
BITMAP KEY ITERATION						
TABLE ACCESS BY INDEX ROWID	DW_LOCATION	1	46	2		
BITMAP CONVERSION TO ROWIDS						
BITMAP AND						
BITMAP INDEX SINGLE VALUE	DW_LOCATION_B03					
BITMAP INDEX SINGLE VALUE	DW_LOCATION_B41					
BITMAP INDEX RANGE SCAN	DW_ORDER_B2					
BITMAP MERGE						
BITMAP KEY ITERATION						
TABLE ACCESS BY INDEX ROWID	DW_PRODUCT	17	1K	10		
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Operation	Name	Rows	Bytes	Cost	TQ	IN-OUT	PQ Distrib	Pstart	Pstop
SELECT STATEMENT		 1	154	34					
SORT GROUP BY		1	154	34	2,03	P->S	QC (RANDOM)	ĺ	ĺ
SORT GROUP BY		1	154	34	2,02	P->P	HASH	l l	
HASH JOIN		1	154	29	2,02	PCWP		l l	
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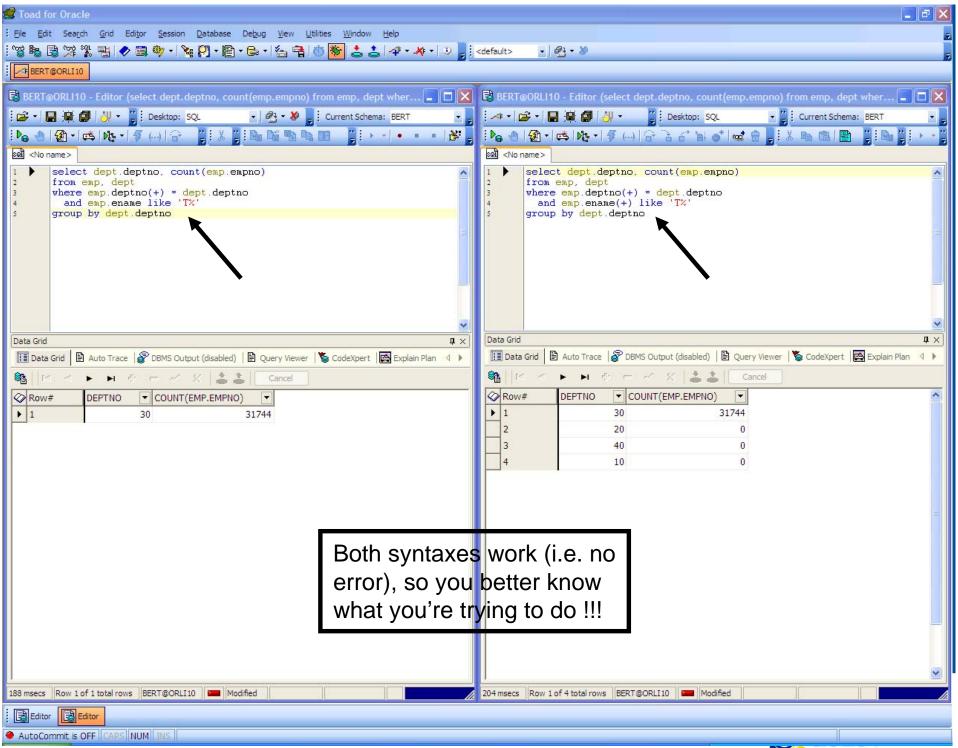
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SQL Guidelines

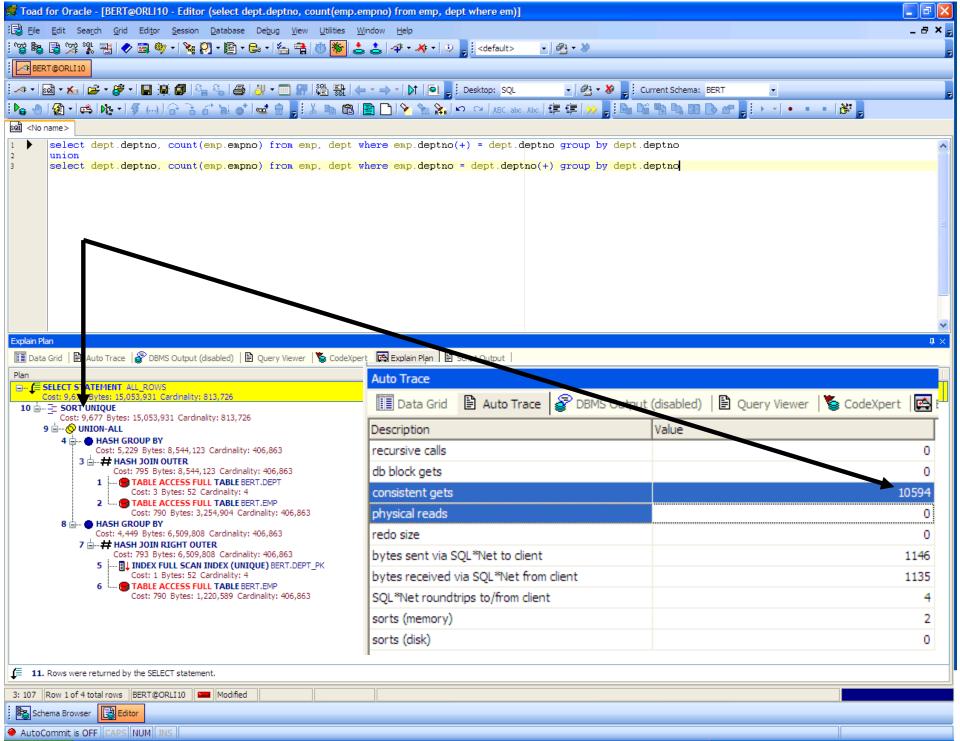
Rule #12: Use ANSI 99 JOIN Syntax – ALWAYS !!!

- •Oracle proprietary (+) syntax has problems:
 - •Cannot do a FULL JOIN efficiently
 - •See slides that follow the next
 - •Outer JOIN syntax prone to user error
 - •You must specify (+) in the WHERE clause for both
 - •The JOIN condition(s)
 - •All other references to that table (source of many mistakes)

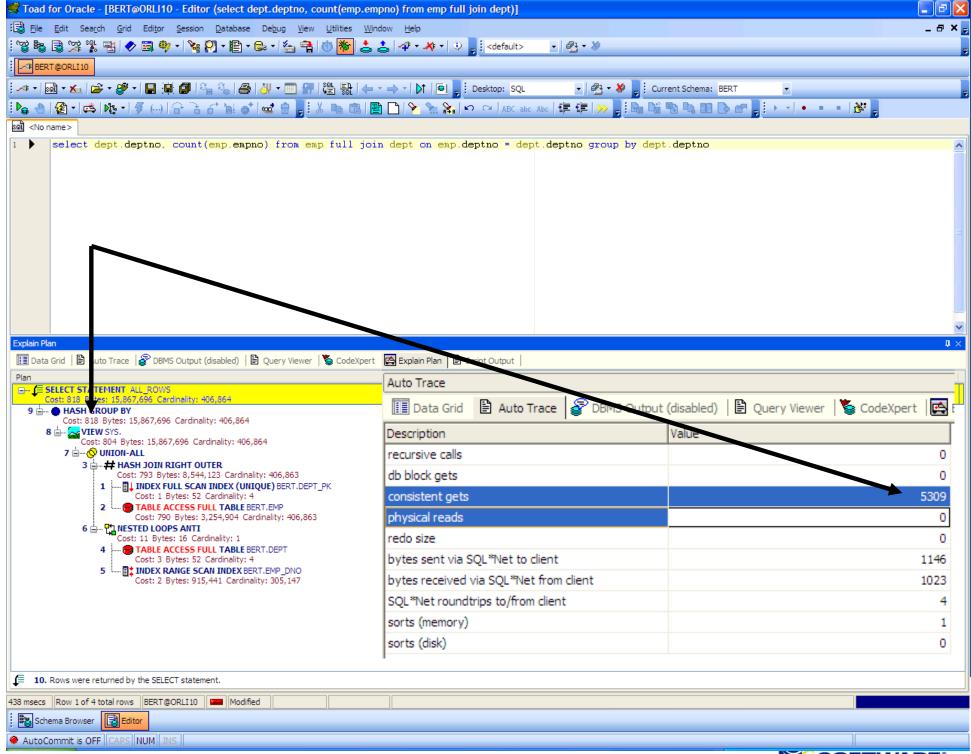












Wow – this is becoming overwhelming

I could go on and list probably another two dozen or so "Best Practices" SQL Tuning and Optimization rules, but you should already be seeing my point – there is a lot of tuning stuff to remember while trying to get your job done.

You should focus on being <u>effective</u> – i.e. the SQL does what the business and/or user requirements mandate.

You should let Toad handle making you SQL efficient !!!

SQL Optimizer knows all this and much, much more: developers can press just two buttons to get their SQL statements automatically and 100% fully tuned!

